Project: DOT The Digital Optical Transceiver Project.

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Audio Samples of Optical Voice





The simplest possible means of communications between two pieces of electronics gear is to run a simple cable between endpoints.



By default, neither *node* expresses any desire to send anything, so they just sit and listen for activity on "the line." Since neither node drives any signal on the line, it tends to float naturally towards some well-known voltage.



Let's pretend our line's quiescent voltage is ground. Suppose A wants to signal to B that some kind of event happened. It may do so by asserting a positive signal on the line. B can readily detect this, because it knows that a grounded wire implies nothing's happening.

Or, is it?



Suppose that node A monitors for one of two events. How can node A inform node B which event occured? We can bring the line high for the first event, but we can't just "assert" 0V on the line to indicate the second event. How can we work around this?

It turns out there are several different ways.



If node B maintains *relative* timing information, it can note *when* node A asserted its signal, and when it stops. By measuring the span of time between these two events, node B can infer whether node A is attempting to indicate event 1 (1 second) or event 2 (2 seconds).

This is called *Pulse Width Modulation*.



If node B maintains *absolute* timing information, it can note *when* node A asserted its signal, and when it stops. By *common agreement*, we know the pulse width is finite, so we instead only care about where the pulse starts in time.

This is called *Pulse Position Modulation*.



While PWM is used with R/C aircraft and cars, it's usually not used in digital communications because it takes different lengths of time to send different numbers. Breaking big numbers into smaller chunks, help, but it doesn't solve the problem completely.



Do we use PPM? In a crude form, we used to use multi-bit PPM for floppy disk recording (FM and MFM encoding). Manchester encoding is a more contemporary application of the idea. Manchester encoding is used in 10-base-2 and 10-base-T Ethernet!

Still, we prefer not to use it, because it basically takes two bits to communicate one.



No, I choose to use PCM -- Pulse Code Modulation. Fancy words for simply choosing to do the simplest possible thing you can: a binary zero and binary one are represented simply by specific voltages (0V and 5V in the case of DOT's hardware).



Which interpretation is correct?!

We need *edges* to keep the receiver in sync with the transmitter, so that the receiver doesn't go too fast or too slow. The trick is inserting these edges in such a manner that we keep our data steam as compact as possible. Three techniques remain in common use today.



One approach is to *scramble* the data using a random number generator. The transmitter and receiver set their RNG to the same initial seed value so they can understand each other.

The disadvantage of scrambling is that malicious users can engineer traffic specifically to counter the effects of scrambling, resulting in loss of synchronization.



Bit-stuffing works through probabilities; since we know when the previous edge occured, we *can predict* where future edges would exist, if they were to occur at all. However, after so many bits, the error introduced from prediction can get so great that you start to misinterpret the signal.

So, we *stuff* bits deliberately, with the intention of enforcing synchronization. Note that stuffing only happens when it's needed!



The final approach is to *convert* 8-bit bytes into 10-bit codewords, each *designed* to have a roughly equal number of 1s and 0s, so as to maintain sufficient numbers of edges that the receiver *never* has more than, say, three 0s or 1s in a row. However, you take a 20%

Arduino Test Bench



How NRZI Works

When we transmit a 1, we do nothing. Otherwise, toggle the output signal.

1 1 1 0 0 0 0 0 0 1 1

Results of Streaming 0 Test

7ED2240600000000000507E 000402642 000000000 7E3923060000000000E37E 000402233 000000000 7ED324060000000000437E 000402643 000000000 7EC924060000000000CA7E # OF 1 BITS 000402633 000000000 7ED124060000000000657E 000402641 000000000 • # OF 0 BITS 7E5A230600000000000547E 000402266 0000000 7EAF24060000000000127E RAW DATA FRAME 000402607 000000000 7ED024060000000000767E 000402640 000000000 7E3F2306000000000897E 000402239 000000000 7ECD24060000000000867E 000402637 000000000 7E4123060000000000FE7E 000402241 000000000 7ECA24060000000000FF7E 000402634 000000000 7ED324060000000000437E 000402643 000000000 7ECE24060000000000B37E 000402638 000000000 7E37230600000000007E(CRC ERROR)

WITH 8-BIT CRC) This is why HDLC

(USES HDLC FRAMING

exists in the first place‼

In case it's hard to read, here's a color-coded transcription of a single data packet from the receiver:

RED is HDLC framing. is number of 0s. YELLOW is number of 1s. BLUE is 8-bit ATM CRC.

7ED12406000000000577E 000402641 00000000

<-- Notice the CRC error caused by the unreliability of the Arduino's RS-232



See how the resulting NRZI-encoded waveform is either almost entirely low voltage or high voltage? This poses a bit of a problem *not* because it's somehow "wrong" to do from a philosophical stand-point. Instead, it pushes the receiver's amplifier to its physical limits, and may actually cause loss of data!

Receiver Schematic (Preliminary)



Receiver Schematic (Preliminary)



Amplifier Saturation Waveform



Thank you for attending! (If Time Permits, Forth Code Follows!)

CRC-8-ATM (sans scrambling)

Polynomial is dreadfully simple: $x^{8}+x^{2}+x+1$

: b dup \$80 and if 2* 7 xor else 2* then ; : c over c@ xor b b b b b b b b swap 1+ swap ; : crc vars 0 c c c c c c c c nip 255 and ;

Proper ATM-spec CRC also XOR's final value with \$AA for scrambling purposes. I'm communicating over RS-232 and USB, so no scrambling necessary.

Primordial HDLC Framing and Byte-Stuffing

: f	Elag	<pre>begin recvRx \$7E = until ;</pre>
: -	-escape	$dup \ \$7D = if$
		drop recvRx \$20 xor then ;
: (D	recvRx -escape over c! 1+ ;
: r	1	0000;
: t	celemetry	<pre>vars n n crc recvRx xor if ." (CRC ERROR)" then drop ;</pre>
: 1	rxFrame	flag telemetry flag cr ;

Notice that CRC byte covers message data *after* HDLC escaping has occurred. Real HDLC wouldn't do this, but it took less code to make it work this way, and it works fine.

Complete Telemetry Source Code

: deviceName

S" /dev/ttyUSB1" ;

variable hRx	
: openRx	deviceName r/w bin open-file throw hRx ! ;
: closeRx	hRx @ close-file throw ;
: sendRx	hRx @ write-file throw ;
variable buf	
: recvRx	buf 1 hRx @ read-file throw drop buf @ 255 and dup hex
	s>d <# # # #> type decimal ;
And the second second	
create msg 0 ,	
: askRx	msg 1 sendRx ;
create vars	2 cells allot
	constant nObits
	constant n1bits
: flag	begin recvRx \$7E = until ;
: -escape	dup $\$7D = if drop recvRx \$20 xor then ;$
: 0	recvRx -escape over c! 1+ ;
: n	0000;
: b	dup \$80 and if 2* 7 xor else 2* then ;
: C	over c@ xor b b b b b b b swap 1+ swap ;
: crc	vars 0 c c c c c c c nip 255 and ;
: telemetry	vars n n crc recvRx xor if ." (CRC ERROR)" then drop ;
: rxFrame	flag telemetry flag cr ;
: .cols	." NOBits N1Bits" cr
	."" cr ;
: .cell	s>d <# # # # # # # # # # # > type space ;
: .row	nObits @ .cell n1bits @ .cell cr ;
: run	cr .cols begin askRx rxFrame .row 1000 ms again ;