

FEFFF

Forth Engine Factored For FPGA

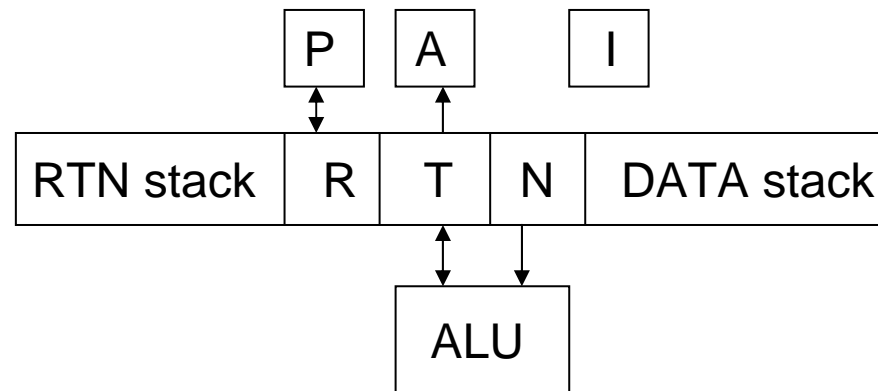
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SVFIG – 7/28/2007

Why? ... Why Not?

- Always wanted to implement my own CPU
- Had some applications needing several orders of magnitude speedup
- Profiled application inner loops, range of architectures... realized core was essentially P16



Current FPGAs

- Have large on-chip memory, roughly 64 K BYTES!
- Register/stack access time about the same as block memory
 - 8-9 ns for 16 x 16 bits of local memory
 - 10 ns for 1K x 16 bit dual-ported block memory

So Let's Re-Factor the P16

- On-chip memory *only*
- Separate data and program memories
 - Not really separate – dual ported block memory
 - Busses rather than muxes
- KISS
 - Not pre-fetching several instructions in one word
 - Hardware blocks idle –vs– pipeline and other complications

Revised Instruction Set

- Execute all opcodes in 2 clock cycles with “simple” clocking
- Several P16 operations broken into two opcodes:
 - Proceed LD, LDP, LDA, POP with a DUP
 - Follow ST, STP, STA, PUSH with a DROP
- PASS
 - TOS \leftarrow Tbuss
- OVER conveniently omitted

Problems with LIT

- P16 grabs next (program) memory word for LIT
- In my implementation, critical path is instruction fetch and decode
- Would have had to latch instruction (and other complications) slowing *all* instructions by some 30%
- So, like JMP, JZ, JNZ, literal is embedded in the instruction word and becomes
 - DUP (preserve current TOS)
 - LITL (load low-byte of literal into TOS)
 - LITH (load high-byte of literal into TOS)

26 Opcodes

JMP to 10-bit address

JZ

JNC

Spare

RET

CALL

Spare

Spare

ADD

SUB

COM

XOR

AND

SHR

SHL (multiply)

PASS (no ADDC)

'LD

'LDP

ST'

STP'

LITL

LITH

Spare

Spare

'LDA

STA'

'POP

PUSH'

DROP

DUP

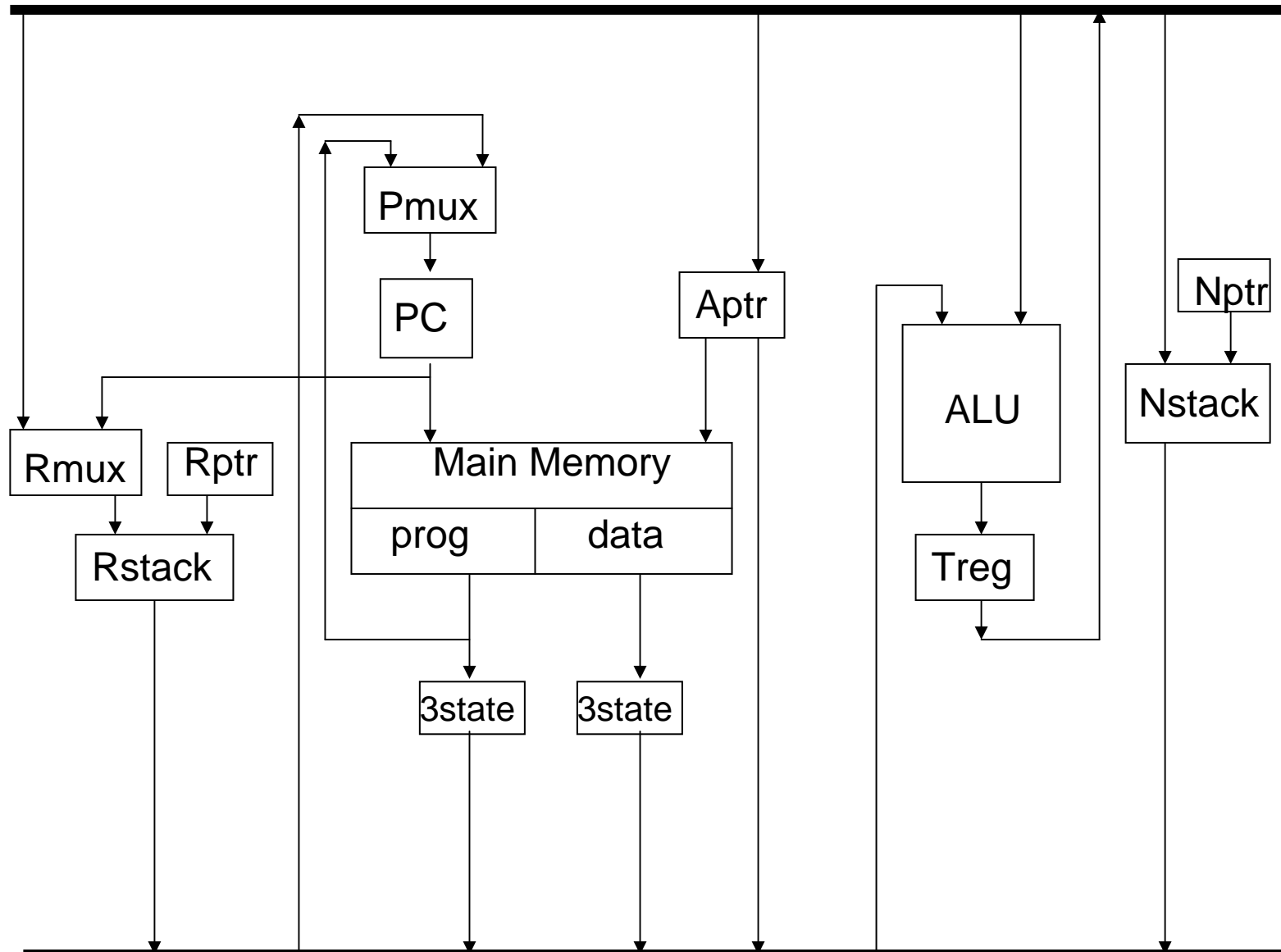
(OVER)

NOP

The Implementation

- Targeting Spartan XC3S1000
 - 48 KByte on-chip memory
 - Future multiprocessing
- 16-bit data and program words
- Return and data stacks 16 words each
- No TMUX
 - PASS opcode sends Tbus through ALU unchanged

Tbuss



Nbuss

Putting It Together

- 8 pages VHDL + 1 page schematic
- Key data paths
 - $N_{stack} + N_{ptr} = 16 \text{ nS}$
 - $R_{stack} + R_{ptr} + R_{mux} = 20/25 \text{ nS}$
 - $ALU + T_{reg} = 17/20 \text{ nS}$
 - $\text{Data memory} + A_{ptr} = 23 \text{ nS}$
 - $\text{Program memory} + P_{mux} + PC = 25 \text{ nS}$
 - $\text{Program memory} + \text{Instruction decode} = 25 \text{ nS}$

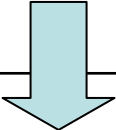
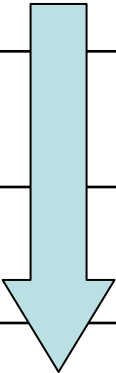
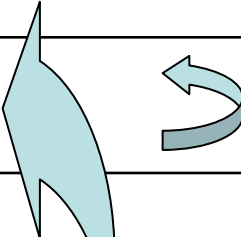
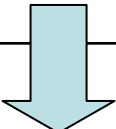
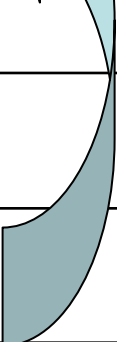
Design Comments

- Xilinx free tools surprisingly good
 - Good synthesis from behavioral description for data path
 - limited simulation time (1000 ns), but good enough
- The “big burden” is the controller block
- Ting’s “bootstrapped debug” *very* helpful
 - Test JMP, NOP, LIT 1st – you’ll need ‘em for other tests
 - Then test ADD and JZ, then DUP and JNC
 - Then CALL and RET
 - End tests with a opcode that jumps to itself – which one (address) well tell you results of conditionals
- Simulation + full-speed hardware “good enough”
 - Hardware single-step and VGA “monitor” very useful

Results & Status

- “Works” in simulation
 - 40 MHz, \leq 20 MIPS
 - Ignoring main memory, CPU uses 3% of chip resources!
- Works in real-world!
 - Tested at 25 MHz
 - Single-step and VGA character generator added
 - 4% of chip logic
 - 8% of chip memory

Single-Step Demo

| Loc | Hex | Instruction | |
|-----|------|-------------|--|
| 0 | FC00 | NOP | |
| 1 | 2804 | CALL 4 |  |
| 2 | 0002 | JMP 2 |   |
| 3 | 0003 | JMP 3 | |
| 4 | 6755 | LITL 55 | |
| 5 | 2000 | RET |   |