# Simple <br> File Query 

## A Forth Standard?

## Dual-CFA Definitions

## Batcher's Sort

## Getting Started with F83

## Windows for the TI 99/4A

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## Forth Dimensions

Published by the
Forth Interest Group
Volume VIII, Number 4
November/December 1986
Editor
Marlin Ouverson
Advertising Manager Kent Safford Production
Cynthia Lawson Berglund Typesetting LARC Computing
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Subscription to Forth Dimensions is free with membership in the Forth Interest Group at $\$ 30$ per year ( $\$ 43$ foreign air). For membership, change of address and to submit items for publication, the address is: Forth Interest Group, P.O. Box 8231, San Jose, California 95155. Administrative offices and advertising sales: 408-277-0668.

## Symbol Table



Simple; introductory tutorials and simple applications of Forth.

Intermediate; articles and code for more complex applications, and tutorials on generally difficult topics.


Advanced; requiring study and a thorough understanding of Forth.


Code and examples conform to Forth-83 standard.

Code and examples conform to Forth-79 standard.


Code and examples conform to fig-FORTH.


Deals with new proposals and modifications to standard Forth systems.


FEATURES

## 13 Dual-CFA Definitions, Part Two

by Mike Elola
The dual-CFA structure provides a new method for decomposing functions into smaller functions. Its value can be demonstrated in deferred and vectored definitions, and in definitions that dispatch multiple functions. This strategy can be the basis of a Forth programming philosophy aimed at compactness, brevity and programming ease.
17 Simple File Query by Edward Petsche


This program allows the user to define and initialize a file, enter data, query on any combination of fields, delete records and change field values in records. It is based on data-base elements presented previously in Forth Dimensions and should work with most versions of Forth-83.
28 A Forth Standard?
by Glen B. Haydon
Forth does not differ from a natural language: it is evolving. And what is a standard language? Only after a word is used with a specific meaning for some time do dictionary editors accept it. This essay considers common use as a commonsense paradigm for Forth standards.
34 Windows for the TI 99/4A

## by Blair MacDermid



This program plots algebraic functions in a choice of five windows on the display. It computes the coordinates of a plotted function, appropriately scaled to fit within the selected window. (Members of the Fort Wayne FIG Chapter implemented the ACM SIGGRAPH CORE Standard as a group project, from which this code was adapted later for publication.)

## 37 Getting Started with F83

## by Greg McCall

Sifting through F83's source shadow screens can be a bewildering first exposure to that system. This summary of the file words and file-editing facilities will ease your introduction. It explains how to open a second, read-only file and load screens from it without changing the CURRENT file.

## 39 Batcher's Sort

## by John Konopka

An alternative to the sometimes quirky Quicksort was discovered by K.E. Batcher - slightly slower, but more robust and with consistent sorting times. If you'd rather not complicate your Quicksort code to handle special cases, Batcher's may be just the sort for you.

## DEPARTMENTS

5 Letters<br>12 Editorial: "Conventions"<br>27 Crossword<br>31 Advertisers Index<br>42 FIG Chapters

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[^0]
## Fast SEARCH for F83

## Dear FIG,

I am happy to finally contribute something to the Forth community. For all of the 8086/8088 F83 users out there, here is a SEARCH function completely written in low-level code that executes very quickly. Since the original F83 SEARCH function was threaded code it was tolerably slow, but a project I've been working on lately needed a quicker search, so I bit the bullet and did it. Here, the function is adapted to the Laxen \& Perry system ... enjoy faster searching!
(In order to maintain the threaded code "purity" of the UTILITY.BLK file, this function should be placed in either the KERNEL.BLK or the CPU8086. BLK source files, and the existing search function in the UTILITY.BLK file should be commented out.)
I have been programming exclusively in Forth for the past three years and, having written both $\mathrm{Z80}$ - and 8086-based systems, I feel qualified to say that Forth offers the greatest man/machine interface yet devised in software. Although it is slightly more difficult to adapt to Forth's subtle programming philosophy, the rewards are quick in coming. I know of many things that can be done in Forth but which are impossible in other programming languages.

As a rather lazy person, I would like to commend all of the FIG community for their tireless efforts in promoting the very best programming language yet designed. And special thanks to Chuck Moore, Leo Brodie, Henry Laxen, Michael Perry, Marlin Ouverson, Bill Ragsdale and all of the other regular contributors to the progress of Forth Dimensions. I hope their example motivates more people to contribute.

## Sincerely,

## Bill Zimmerly

St. Charles, Missouri

## Natural Word Usage

## Dear Mr. Ouverson:

Ting's computation of static F83 word reference counts ${ }^{1}$ is the first I have seen. The total number of words $(11,063)$ is large enough to be interesting. I immediately plotted a graph with the words ordered by frequency of use. A $\log -\log$ plot was the cleanest and had, for me, a surprising result: reference count was inversely proportional to frequency, i.e., the data closely fits a line of slope -1. I tried several other populations I had available ${ }^{2,3}$ : one of spoken English and one of written English. The results were the same!

While browsing at the library one day, I came across a volume on Zipf's
law ${ }^{4}$. The explanation was at hand: this is a property of human behavior. Thus, Forth has some of the properties of natural languages.

I also investigated various coding techniques ${ }^{6}$ to determine the amount of compaction that can be obtained taking advantage of the frequency-ofuse statistics. The results are somewhat disappointing. For hardware implementations, a block encoding is probably all that can be justified.

| Number of different tokens (words) | 555 |
| :--- | ---: |
| Total number of occurrences | 11,063 |
| Block code size | 9.116 bits |
| Theoretical code size | 7.051 |
| Hoffman code size | 7.084 |
| $4-8-12$ repeated comma code | 7.821 |
| $4-8-12$ non-repeated comma code | 7.735 |
| $8-16$ repeated comma code | 9.316 |
| $8-16$ non-repeated comma code | 9.316 |

In the repeated codes, the same token (word) can be coded in several sizes which, of course, lowers the coding efficiency. The relatively small number of words (compared to $2^{* *} 16$ ) accounts for the poor performance for the 8-16 codes.

1. C.H. Ting. "F83 Word Usage Statistics." Forth Dimensions VII/4, pg14, November/December 1985.
2. H.F. Gaines. Cryptanalysis. Dover, 1956.

## Zimmerly's F83 SEARCH

## B6 <br> \String functions... <br> ASSEMELER Label (FIND1) <br> DX SI MOV BX DX MOV BF FOP $2 P U S H$ <br> CODE SEARCH ( SADR SLEN BADA BLEN -- NF) | FIND SUPGTfing

CLD CX POF DI FOF EX FOF DX FOF EF Pugh DK 51 XCHG
CS AX MOV à ES MOV ij [SI] AL MOV here gyte ref SCAS 0= If EX FUSH SI FUSH DI FUSH DI DEC EX IX MON ( get count GYTE REP 2 CHPS $0=$ ( cocapare the strings for equality!) If EX FOF AX FOF AX PDP BX DEC - $\#$ \# AX NOV : true flay! (FINDI) H) JMF THEN DI POP SI FOP CY PMF ELSE AX AK YOR
(FINDI: \#) JMF THEN \#) JMF END-CODE

```
223
```

223
String functiong... WBT 11-09-1085
String functiong... WBT 11-09-1085
(FINDU! is the exit point for the SEARCH furition that follows.
(FINDU! is the exit point for the SEARCH furition that follows.
At entry, the 暗contalis the offset address, and the
At entry, the 暗contalis the offset address, and the
AX contains the TRUE or FALGE flag.
AX contains the TRUE or FALGE flag.
SEARCH is a very high speg thow could it be faster?) function
SEARCH is a very high speg thow could it be faster?) function
that scans a string trying to lorate the glven autstring
that scans a string trying to lorate the glven autstring
within it. The method used is to searct for the first
within it. The method used is to searct for the first
chafdcter, and when found, compare the characters that
chafdcter, and when found, compare the characters that
follon it for a complote match. If both fail, the sear=h
follon it for a complote match. If both fail, the sear=h
for the first character continues trom where it left gif
for the first character continues trom where it left gif
until we've scanned the entire buffer.

```
        until we've scanned the entire buffer.
```


3. G.D.A. Brown. A frequency count of 190,000 words in the LondonLund Corpus of English Conversation. Behavior Research Methods, Instruments \& Computers, 16 (6):502-532, 1984.
4. S.R. Ellis and R.J. Hitchcock. 'The Emergence of Zipf's Law: Spontaneous Encoding Optimization by Users of a Command Language.' IEEE Transactions on Systems, Man, and Cybernetics, vol. SMC-16(3):423, May 1986.
5. G.K. Zipf. Human Behavior and the Principle of Least Effort. Addison-Wesley Press, 1949.
6. R.W. Hamming. Coding and Information Theory. Prentice-Hall, 1980.

Sincerely,
James C. Brakefield, M.S.E.E. San Antonio, Texas

## Leaky Sieve

## Mr. Ouverson:

In the process of optimizing the sieve benchmark, Terry Noyes has unwittingly rejected a superior algorithm and departed from the de facto benchmark standard. The sieves Mr. Noyes calls 'corrupt" are not corrupt at all. They correctly count 1899 primes from 3 to 16383 . The FLaGS array represents only odd integers, and only odd multiples of primes are "flicked." Fortunately, the Noyes version is easily modified to incorporate the better algorithm, and the resulting version finds the 1899 primes slightly faster than the unmodified version finds 1028.

## Stephen Brault

Chandler, Arizona

## Mr, Ouverson:

I would like to retract my April letter to you (Forth Dimensions VIII/2) and live in shame for the rest of my life.

I had thought that the $0-8192$ loop in all sieve benchmarks meant they were calculating the number of primes from zero to 8192. Not so. A few weeks after the letter was sent to you, someone pointed out that these sieves were actually finding primes in the range of zero to $16,000+$ by looping through the 8192 odd numbers in that range.

Oh.
Fortunately, we use identical code to benchmark other Forth systems, so they also received the five percent speed improvement resulting from using the wrong sieve program. I've enclosed the proper Forth sieve with this letter.

Living and Learning,

Terry Noyes<br>Palo Alto Shipping Company<br>Menlo Park, California

## Seeing is Believing

Dear Marlin:
I enjoyed Michael Ham's "Making Numbers Pretty" (VII/5). I had just written a routine to calculate the necessary number for masking a given bit (or bits) and Michael's words .BITs and 16BITs fit in perfectly, although I modified them slightly.

Referring to the enclosed listing, the words BIT-MASK and 2BIT-MASK return to the console the number necessary to mask the desired bit, or bits, in the current base. The words .BITs and 16BITs visually confirm the mask, making life a little easier for us doubters!

For example, if you want to mask bit five, then entering 5 BIT-MASK (in base ten) will give:

32 In base 10
HI: $0000000000100000: L O$
while if bits four and six are to be masked, then 46 2BIT-MASK (in base sixteen) will return:

50 In base 16
HI: 0000000001010000 :LO

```
Noyes' Sieve
decimal
8192 constant size
variable flags size vallot
primes ( - primes) < does the primes once)
    flags size 01 fill { initlalize the array)
    0 (prime counter)
    size ( range/2 of numbers to do)
    flags i + ce (see if prime already)
            IF
            3i+i+dup i + size<<<don't go too far)
                size flags + over flags + i + (range of nums to tag)
                a i c! dup (tag numbers as non-primes)
                +LOOF
            THEN
            drop 1+ (drop the i used for +loop, increment prime count)
            THEN
    LOOP;
: sieve
    COUNYTER < start counting )
    10 0 DO primes LOOF (perform10 iterations)
    TIMER (stop counting )
    CR.." primes" (print the number of primes )
        90 DO DROP LOOP; (Elean-up stack)
CR . (Type 'sieve' to execute this benchmark program ) CR
```

Forth Sieve. Uses pointer arithmetic to calculate the number of primes from zero through $\mathbf{1 6 3 8 3}$. To save space and time, it only needs to work with the 8192 odd numbers.

```
Thomas' Bit-Mask Locator
Listing d
serean 脑
    I masking-number calculator gtAugge
    8PC 32 HOLD I
    | 16BITS ( FD 7/5, M.Ham, modified) CR " HI|"
```




```
    : EITS (FD 7/5, M.Ham, modifimd) EASE SWAP
        2 BABE ! 8->D 16BITE BAEE ! QUIT ;
    GAS= BASE DUP DECIMAL " In base " , EABE ! I
    %
    , BEE-MAgK DUP CR U. BAB= .BITS I
    MABK DUP IF 1 SWAP ELA (lift shift) EL8E I OR THEN I
    BIT-MAEK ( n -- ldisplay numbar to mask bit ni n=0 thru 15)
        MASK BEE;MASK I
    2BIT-MASK ( nl n2 -- ldisplay numbur to mask bits n1 & n2)
        MABK SWAP MASK + SEE-MASK I
```

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Conveniently, both the base of the mask and its binary representation are displayed. (Remember, the sixteen bits are numbered zero through fifteen.)

The word SLA in MASK is my system's ML shift-left arithmetic word (n1 cnt -- n2). Replace it with your appropriate instruction. The 1 OR in MASK takes care of the zero bit position, as in 0 BIT-MASK.

Forth Dimensions and its contributors often supply me with either some finishing touches or an idea to expand on. Thanks!

Sincerely,
Gene Thomas
Little Rock, Arkansas

## Student Roots

Dear Editor,
During this Summer Quarter of 1986, I have been providing the coursework for a student taking "Forth Programming" at Auburn University at Montgomery. As one of his assignments, this student (Hunter Moseley) was required to write a square root in Forth (F83) based upon a Newton's method-type algorithm. However, Hunter went beyond my thought and wrote code that put mine to shame. My code is shown in Figure One.

The $\mathrm{D}^{*}$ I used does the same thing as */ but with double-precision numbers. In other words, (d1 d2 d3-- d4). Also, the 2NIP is a double-precision NIP. I hated to use the double-precision words, but for the accuracy needed, they were necessary.

Hunter's code was simply that shown in Figure Two.

In a time test on a Zenith- 151 with 10,000 iterations, dropping the result each time, Hunter's code guaranteed 119 seconds with any input from zero to 32,766 . Mine, however, with an equivalent range of inputs, does the square root of one in seventy-five seconds, the square root of two in 280 seconds, and gets even worse after that.

As can be seen, the two approaches are based on the same idea, but Hunter's does no bound checking. His

simpler application of the algorithm is much slicker - beauty in Forth.

Additionally, as an experiment with F83's assembler, I translated Hunter's algorithm into assembly. The code is listed in Figure Three. A time test on the Zenith-151 with 10,000 iterations, dropping the result each time, guaranteed five seconds! Yes, that's right 2,000 iterations per second! Perhaps this amazes no one else, but I was somewhat shocked.
For those interested, Hunter also has the signed, double-precision version of the square root. The code is in Figure

Four. The $\mathbf{D} /$ is a double-precision divide. If anyone is interested in the code for these operators and their double-precision primitives, I will gladly share them.

In any case, I present these attempts as examples of how traditional mathematical thought sometimes must give way to the more efficient patterns used by our friends - the computers - and Forth.
Sincerely yours,
R.L. Davies

Montgomery, Alabama

## Second Take: <br> Multiple LEAVEs by Relay

## Dear Mr. Ouverson:

Please discard my previous letter to you (Forth Dimensions VIII/3), as it was completely erroneous. My intended verification test wound up with confusion between the fig-FORTH words in my system and the new words, due to my carelessness! Here is the new manuscript:

John Hayes' "Another Forth-83 LEAVE" (VII/1) stimulated me to try to find an even simpler way to handle multiple Forth- 83 Leaves. I decided that a straight-forward approach involved having each LEAVE simply branch to the next LEAVE, with the last one removing the index values from the return stack and branching to the word following LOOP.

I 'grafted"' such a construction onto fig-FORTH with the definitions below; words with a * prefix are used to identify changes from fig-FORTH. Unstarred words such as (DO) and (LOOP) are unchanged. Whenever a *LEAVE is compiled, the variable place is used to hold the location of its branch value for later adjustment. This variable also serves as a flag to show that there is a leave branch to be resolved. *LOOP calls a >RESOLVE to install the jump value of the preceding (if any) *LEAVE; also, if there is a *LEAVE in the word, a special OUTLEAVE is compiled immediately following the (LOOP) closure. OUTLEAVE removes the (two) loop parameters from the return stack and proceeds to the next word, i.e., the word that was entered after *LOOP. If the *LEAVE command is not invoked at run time, the normal loop operation removes these parameters from the return stack, so OUTLEAVE must be skipped over. *LOOP compiles this bypass with a BRANCH 4 which is encountered in normal loop completion. Alternatively, (LOOP) could be modified to use OUTLEAVE in normal loop completion.

Note that oUTLEAVE can be a primitive which removes two words from the return stack by using PLA four times. If outleave is defined as a

```
Page's LEAVE by Relay
] UARIABLEE FLACE
        FESOLVE HEFE RUEF - SWAF : ;
        FESILUE HERE - :
    IUTLEANE F:F: DFOF FY DROF F :
    - [OD OLACE COMPILE GOS HERE 3 ; IMAEDIATE
    Same as FIg D⿴囗⿱一𫝀口儿
```



```
        PLAEE A FDUF IF DOMAFILE RRANEH A , Y for gKIFP:MIG DITLEAUE
        FESOLUE COMFILE DUTLEAME THEN O PLACE: ; IMMEDIATE
    *LEME FLAEE a OQHF IF >REGOLUE THEN S resolue zny pregeding LEGUE
    COMPILE BFANGH HERE FLACE ! 0, : IMMEDIATE
```



```
    I 10 = IF I . -LEA''E THEN *LOOF ." ENL " ;
    TEST1 5 O POO = 1 *DG J S = IF ! . *LEAUE THEN *LODP ." END " HOQF :
: TEST2 5 g *DO 6 1 *DO *LOOF I 3 = IF I . *LEMUE THEN *LOOP ." END " ;
    TEST: ח1 rz--- * OD 7 0DO I 3 =
    IF l. *LEANE THEN *LOOR." INNEER"
    15= IF 1 , *LEAUE THEN *LODF ." END " ;
```


## Borenstein＇s Fixed－Point Trig

```
SCR * 1
0(SIN SCALED BY 3784)
1 (0< X<59444 = 0< X<90 DEG)
2(0<SIN X<3784 < 0<SIN X<1)
3 O VARIABLE XS
    DEG 6604 100 */
    KTIMES 17321 M* SWAP DROP ;
    TERMI XS @ U* DROP U* SWAP DROP MINUS ;
    SINI
    DUP 256 > IF }\quad\begin{array}{cl}{\mathrm{ DUP (Check on small X )}}\\{\mathrm{ (Leave one copy of X on stack)}}
                DUP U* SWAP DROP DUP XS! (XS=Z*Z/2*16)
                4U* DROPMINUS (2n16-4*XS)
                    7 TERMI 15 TERMI 50 TERMI
            U* SWAP DROP
        THEN;
SCR*2
    (COSI TANI SIN COS TAN)
    ?MIRROR DUP 5944 > IF 11888 SWAP - THEN;
    REDUCE 23776 MOD DUP O<IF 23776 + THEN DUP 11888<IF
        ?MIRROR ELSE 11888- ?MIRROR MINUS THEN ;
    SIN REDUCE DUP ABS SINI SWAP O< IF MINUS THEN;
    COSI 5944 SWAP-SINI;
    COS 5944 SWAP-SIN;
    TAN1 DUP SINI 3784U* ROT COS 1 461 MAX U/MOD SWAP DROP ;
    TAN 11888MODDUP 0<IF 11888 + THEN DUP 5944 > IF
        11888 SWAP - TANI MINUS ELSE TAN 1 THEN ;
```

colon word，its first operation is to call the generic colon－word procedure which pushes the compilation address of the following word（the desired one at the end of the loop）onto the return stack，above the parameter values to be removed．In this case，the top value must be saved and restored by defining：
：OUTLEAVE
$R>R>D R O P$
$R>D R O P>R ;$

The test words use this＊Leave in multiple occurrence in a single loop；in single occurrence in both inner and outer of two nested loops；and in both inner and outer loops．In nested loops， any＊LEAVE in an outer loop must occur subsequent to the end of the inner loop．Starting a new loop before the forward resolution of the＊LEAVE jump would cancel the record of the ＊LEAVE（in PLACE）．

Chester H．Page
Silver Spring，Maryland

## MMS Upgrade Offer

## Dear Editor：

Any licensed MMSFORTH user who had not received a gold－colored v2．4 discount letter by the end of September 1986 －please notify MMS of your current address so we can send it along， or call us for further information．

Sincerely，
A．Richard Miller
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## Fixed－Point Trig

In the May／June 1986 issue of Forth Dimensions，we carried an article titled ＂Fast Fixed－Point Trig＇’ by Johann Borenstein．Due to space limitations， we were unable to print the companion screens to the article in that issue．You will find them herewith．
－Editor

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## Conventions

We've talked about this before, but someone - one of our authors, yet recently confused F83 with Forth-83. Big mistake! 883 is an ultra-superset of Forth-83, nearly an order of magnitude larger. Forth-83 and 83-Standard are common shorthand for the phrase, "Forth-83 Standard." The name F83 is not an even shorterhand, it is the name of an implementation of Forth. Does everyone understand the difference between a language implementation and a standard?

As we send this issue to press, last minute preparations are underway for the imminent 1986 Forth National Convention. A major component of this year's convention is a six-part seminar on Forth engines. Those sessions will focus on the new multiplestack WISC (writeable instruction set computer) machine; applications of the Novix 4000 and the design of the Novix 6000 chip; Forth engines developed by Hartronix, Lockheed and Johns Hopkins; ROM-based Forth engines (i.e.,
the Super-8, R65F11 and F 68 HCl 1 microprocessors); Forth engine software; and the future of Forth engines.
Numerous additional events and concurrent sessions are planned to serve the particular interests of all attendees. Exhibitors will include major vendors of commercial Forth hardware and software. Special groups will gather to discuss F83, MacForth/MultiForth, MVPFORTH, NC4000, polyFORTH and 68000 machines (e.g., Macintosh, Atari, Amiga). There will be tutorials about control structure extensions, files and string I/O, multi-tasking in polyFORTH, oblique flying wings, target compilation in F83 and vectored execution of I/O words.
A FIGGRAPH session will feature the latest in computer-generated graphics of significance to the Forth community. FIG chapters' representatives will convene, and there will be a national meeting of Forth Interest Group members. As in past years, the convention will also feature a banquet with
keynote speaker (separate registration required to reserve a seat), a report from FORML including this year's trip to important Forth sites in China, and a "fireside chat" with Mr. Charles Moore, original developer of Forth.

There you have it in a nutshell, or perhaps in a kernel. It seems that as Forth has matured, it has gathered a potency which can propel it into new areas. This integral vitality can lead Forth in unexpected, surprising directions. Keep abreast by joining us in California on November 21-22 at the Santa Clara Trade \& Convention Center, near the new Doubletree Hotel And for an intensive immersion in Forth methodology and experimental proposals, stay for the following weekend's FORML conference at the Asilomar conference grounds in Pacific Grove, adjacent to Monterey. Information for either event can be obtained by calling 408/277-0668, the FIG hotline.
-Marlin Ouverson Editor

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## Part Two：New Solutions to Old Problems

# Dual－CFA Definitions 



Mike Elola<br>San Jose，California

The introduction of a dual－CFA definition structure provides a new building block for Forth programming． It also provides a new method for decomposing larger functions into smaller functions．

Because of the new possibilities af－ forded by dual－CFA decomposition， worthwhile changes to many imple－ mentations of Forth can be formu－ lated．By contrasting these new defin－ itions with the alternatives currently available，the value of this new metho－ dology can be demonstrated．

Changes are suggested within each of three different areas common to most implementations of Forth．The areas to be covered are（1）deferred definitions，（2）vectored definitions and（3）definitions that dispatch multi－ ple functions．This roughly parallels the organization of topics in my orig－ inal paper describing dual－CFA de－ composition ${ }^{1}$ ．It showed that dual－ CFA definitions help maintain a con－ sistent strategy for decomposition and that this strategy can be the basis of a Forth programming philosophy aimed at memory compactness，brevity of expression and ease of programming．

## Deferred Definitions

Deferred definitions are used to al－ low a lower－level word to dispatch a function that is defined in terms of many high－level support words．When definitions that require the undefined function are compiled，a superficial， ＂stand－in＂definition is compiled in lieu of the actual，desired function． Later，the stand－in definition＇s body is modified to reference the correct，high－ level definition．
Dual－CFA definitions can be used to implement deferred definitions．The dual－CFA word functions as the stand－ in definition that is modified later when the dictionary contains the sup－ port needed for compiling the＂real＂ definition．
In this implementation，the dual－CFA definition performs a self－modification
function（aided by the dual－CFA defin－ er）．The child transforms itself into a parentless，single－CFA definition when executed．After execution，it contains a reference to a headerless definition at the top of the dictionary（see Figure One）．

The definition for the parent definer is：


Dictionary entry before execution：


Dictionary entry after execution：
Self－transformation of a dual－CFA word created with DEFER

Figure One

```
: FAILINGG-NUUNEEF: i -- j
    FEEATE DIELOUL, LUAFILE-DEF
    UUES` i síring-adi ciad -- a ;
```



```
    ##
    G O रUT EEGIN
        CUNVERT DLiF CE ぶ - winILE
            yuF CE DUF उB=
            SWMAF 4AG 4E WITHINN +
```




```
        REFEAT \NOUF
    F; 巵; DKUF IF DNEGMTE THEN;
```

Figure Two
The advantages of this technique are the unavailability of the definition－ modifying function except to the unin－ itialized children of DEFER：，the use of only one name field（as compared usually to two）and the decreased chance of crashing（since the definition－modifying action can＇t be applied to just any word）．

Note that the first advantage is ac－ tually a limitation that may not appear advantageous to some．In F83，defer－
red definitions are initialized with is， which patches the specified word．Is can be used more than once and can be used with any word，not just those words that are created with DEFER．

A new disadvantage regarding the dual－CFA implementation is that the deferred function must be specified in a non－standard way：the name of the deferred definition replaces the colon and name string at the start of the definition．Because of this，it is more difficult to separately recompile the high－level definition again．For ex－ ample，DEFER：may be used to create ？ERROR，which is defined at a later time in the following manner：
？ERROR（flag－－）
IF ．＂ERROR＂ABORT THEN ；
To separately recompile ？ERROR again，a colon must be provided．How－ ever，the first body of ？ERROR still points at the original，headerless defin－ ition．
One solution is to provide another support word：

```
:REDEFER: <old deferred word> (-- )
    [COMPILE] , (PFA)
    DEFERRED-CFA SWAP CFA! ;
```

Now you need not change the source code by adding a colon．Instead，you enter REDEFER：PERROR as a preparative step．Then you can load the source code as－is．
In the preceding definition，defer－ Red－CFA is a constant．It points to the DOES $>$ phrase in the parent defining word．Its derivation was not shown． One way to derive it is to use：

## DEFER：JUNK <br> LATEST PFA CFA＠（cfa－value－－） <br> FORGET JUNK <br> CONSTANT DEFERRED－CFA

Note that REDEFER：increases the chance for crashes，since its definition－ modifying function is not restricted to deferred words．To remedy this，extra code can be added to the definition to ensure that it contains an upstream reference：

## MasterFORTH

FORTH－83 ST ANDARD
－． 6809 Systems available for
FLEX disk sustem s ．．．．$\$ 150$
$059 / 6809$ ．．．．．．．．$\$ 150$
－ $680 \times 0$ Systems available for
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CP／M－68K
$\$ 150$
－tFORTH／20 for 68020 Single Board Computer
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EpROM set for complete stand－ alone SBC
$\$ 390$
－Forth Model Library－List handier，spreadsheet，Automatic structure charts ．．．each ．\＄40
－Target compilers ： 6809,6801 ， $6303,680 \times 0,8088,280,6502$

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68020 SBC， $51 / 4 "$ floppy size board with 2MB RAM， $4 \times 64 \mathrm{~K}$ EpROM sockets， 4 RS232 ports， Centronics parallel port，timer， battery backed date／time， interface to $251 / 4$＂floppies and a SASI interface to 2 winchester disks ．．．．．．$\$ 2750$ 68881 ft pt option ．．．．．$\$ 500$ OS9 multitask\＆user OS．．$\$ 350$ FAST！ int．benchmarks speeds are $2 \times$ a VAX780， $10 \times 2 \mathrm{ABM}$ PC

## Listing One

Processing of counted strings（i．e．，already－parsed words）
Target Stack Effecti cadd－－j

```
: FAILING-iUuOkUF: ( <error-processing-funcíion` )
    位位E DUCÜL, CÜfFILE-DEF
    DUES; i cfaz ; 隹 i cadi i pia ien j ilag -- ;
        DuF COUNTEXT E E iFIND) i cadd i pia ien j fiag j
```




```
    NLiHIEER-VALLIET ESTACK i overfion?;
    DFL E Ó IF DNOUF THEN K; F; ZDROF ;
```



```
    N\mp@code{UMBER-VALLUE? DFL E 1+ IF}
```



```
: İNTENFRET-WÜ\overline{\}\overline{y} ( cadi -- ;
    BINTERFRET-NUUHiBEM;
```



```
    ( pia ; CFA EXECLITE TSTACK i underfiow?i ;
: INTEKFFKET-iwüर̄D
    ?CUMFILE-NUMBER;
```



```
    i pia ien ; iğz ; if EFA EXECitE ElSE EFA, EXELUTE THEiv ;
```

Word Parsers

```
Targei Siack Effecii -- fiag;
: iwÜRD; i sirean-adid -- fiag;
    ...TU.EE.SLFFLIED... FiERE CE ;
:TIE-WO゙N\ i c -- fiag;
```



```
; BLKi-wÖRD i c -- +iag;
    ELk E EiLŪGK iwuORDj ;
```

Null－delimited input stream parsers／processors

```
Target Stack Effeci; -- ;
VAM\IABLE FFOOCESS-WUURD'
STREAMi-F\tilde{RUUCESSÜR:}
i iparsing-function-ieaving-counted-string-at-dp;)
        CKEATE DÜLül ; COUNFILE-DEF
```



```
        BEGIN FE EXECLITE WHILE
            HEKE FRUUEESS-WORD' EXELUTE FEFEAT
        \tilde{N}\ \overline{vRUuF ;}
    STREÅi-FROUCESSUR: TIE-FRUCESS i -- ;
        \\overline{L}
```





```
: REDEFER: < old deferred word> ( -- )
    [COMPILE] , (PFA )
    DUP @ U< \(\mathbf{0}=\mathrm{IF}\)
    CR ." - MUST BE A DEFERRED WORD"
    ABORT THEN
    DEFERRED-CFA SWAP CFA ! ;
```


## Vectored Definitions

Vectored definitions can often be replaced by fixed-behavior, dual-CFA words. Such words can directly invoke the desired function. For flexible processing, a variety of these words can be defined. Each one would be suited to use in a particular context.

Dual-CFA words offer flexibility in a fundamental Forth form: compiletime selection of the desired behavior by a reference to the correct word from the dictionary. This practice retains the ease-of-use that characterizes normal, fixed-behavior words ${ }^{1}$.
For example, number often employs a vector to provide a means for flexible processing. Because the behavior you desire is usually known at compile time, you do not really need run-time flexibility - just a wider selection of compilable behaviors. This makes number a good candidate for dualCFA decomposition.

The function of number is to convert an input string into a number. When the conversion process fails, program execution is immediately aborted in many Forth implementations. Such an outcome is fine during the interpret or compile phase, but often is undesirable in a finished application.
Several versions of number are needed. Each would have a different failure outcome. This can be achieved by creating a definer word that incorporates NUMBER (see Figure Two).

To define a number-conversion routine suitable for internal use when interpreting or compiling, one child definition might be:

## FAILING-NUMBER:

## nUMBER-VALUE? (string-addr -- d?) CR COUNT TYPE <br> ." - NOT RECOGNIZED" ABORT ;

To define a version of NUMBER more suitable to an application, another child definition could be:

## FAILING-NUMBER:

INPUT-NUMBER (string-addr -- d?)

## C@ BACKSPACES

TIB 12 EXPECT $0>I N$ !
32 WORD HERE RECURSE ;

The advantages of the dual-CFA definitions over vectored definitions are the memory compactness of the compiled words, the absence of intermediary variables, the absence of required initializations and greater immunity to crashes.
Normally, decomposition of the error-handling code within NUMBER would not be possible, unless such code is moved outside of the begin while repeat loop. This would allow the error instructions to be separately specified within any number of parent definitions.

This is similar to the approach taken in F83. This Forth implementation includes a primitive version of NUMBER that does not abort on error. Instead, it leaves a flag on the stack that can be used by parent words to trigger any kind of error processing desired. Since parent words must process the flag left on the stack, a conditional phrase is normally required in all the parent words where (NUMBER) is used.
 sive royalty arrange-arrangements As your long as your applications Each version is the basic FORTH number of pop system, you are
iree to make as number of po systems
including Including XT and AT for compatibles); the TRS-80
Model I II and $4 / 4 \mathrm{P}$; or any $\begin{array}{rr}\text { FORTH } \\ \text { computer system } & \text { Irom }\end{array}$ computer system $=0$ SOTA 15 the
running $\mathrm{CP} / \mathrm{M}$ (version 2.x)
or CP/M Plus
(version 3 $x$ ) chonce for both whats more What's more,
SOTA doesnt SOTA doesn't
require you require you
to enter into any awkward
 experienced programmer When you copy today model and 79 stand SOTA, both the lig with the following exard come complete with the rollowing extra features at no
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Availability: F 68 HC 11 Production units with Max-FORTH ${ }^{\text {re }}$ in internal ROM available 4Q/86. Volume quantity available 1Q/87. X68HC11 emulator with MaxFORTH ${ }^{\text {m }}$ in external ROM available now. NMIX-0022 68HC11 Development System boards available now: $\$ 290.00$

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To avoid having to repeat that failure processing with each use of the F83 number primitive, enlarged functions could be created. If desired, versions equivalent to INPUT-NUMBER and number-value could be created. Such versions would exhibit the same ease-of-use as the dual-CFA versions.

The dual-CFA versions would retain a very slight advantage over their F83 equivalents: they should compile in less memory and should execute slightly faster due to a decreased number of conditionals.

## Definitions That Dispatch Multiple Functions

When implementing function-dispatching words, dual-CFA definitions can also be advantageous. Listing One includes several examples that help illustrate those advantages.

In most Forth implementations, the main function-dispatching routine is INTERPRET. Both the compiling and interpreting functions are often performed within interrpret. Since there is so much commonality between these two distinct functions, it is easy to think of them as children of the same parent process. But what exactly should this parent process be? The answer can be found by more clearly discerning what functional areas are to be combined.

The common ground between the compiler and the interpreter is the input parser. The input parsing function is the same, whether compiling or interpreting. It remains a static function even if there are mid-line transitions between the compiling and interpreting functions.

But because of input redirection, the input-parsing function is not always static. When loading a block, the input source must be the block buffer, not the text input buffer (TIB). Normally, this flexibility is achieved as a run-time function of word. So word normally has a variable behavior dispatched through a conditional phrase. A majority of the time, this conditionallyvariable behavior can be eliminated. Input redirection is rarely exercised during run time. For those exceptions, a variable-behavior version of WORD can be defined by referencing the fixedbehavior versions. By defining LOAD
with a fixed-behavior version of wORD that only parses blocks, the input redirection required by LOAD is enabled using only fixed-behavior words at compile time.

The implementation shown in Listing One factors the function of inputstream parsing to a parent definer. Each of the two dual-CFA children dispatch a different version of wORD. The correct one can then be selected at compile time to suit a given context. Refer to STREAM-PROCESS, TIB-PROCESS and blK-process in Listing One (as well as PRINTING-LOAD in Listing Two).

To provide additional, run-time flexibility, WORD can be defined in terms of the new primitives:

## : WORD ( $\mathrm{c}-\mathrm{-}$ ) <br> BLK @ IF BLOCK-WCRD <br> ELSE TIB-WORD THEN DROP ;

To make a nicer, error-detecting version, the flag returned by the primitive versions of word could be processed as follows:

## : WORD ( $\mathrm{c}-\mathrm{-}$ )

BLK @ IF BLOCK.WORD
ELSE TIB-WORD
THEN $0=I F$
CR ." UNEXPECTED END-OF-INPUT."
ABORT THEN ;
The variable-behavior version of wORD is needed for single-word parsers such as ' (tick) and Create. Having all three versions (WORD, TIB-WORD and BLOCK-WORD) provides the programmer with more choices. Why use the variable-behavior version of WORD, with its extra overhead, when inputredirection flexibility is not necessary at run time?

As an extensible programming language, Forth can exhibit a wide range of functionality that broadens with each new word added. For every programming problem confronted, Forth can be extended in ways that make the solution easy to program. Not only is the original problem more easily solved, but also many related problems become easier to solve.

See Listing Two for a printing version of LOAD that is defined very simply. It could be useful on those few occasions when a screen will not load
(Continued on page 32.)

# Simple File Query 

Edward Petsche<br>Greenport, New York

This article describes an implementation of a simple file query based on the data-base elements presented in Forth Dimensions (see volumes three and four). The parameter fields of words defined by FILE and FIELD have been extended to include some new parameters necessary for the query. The doer and MAKE vectored execution words described in Thinking Forth are also used in this program. Implementations of these words for various versions of Forth are given in that book's appendix. If you don't have access to that book, the implementation in screen 8 should work for all versions of Forth-83. If you prefer the DEFER and Is vectored execution words, the necessary modifications, aside from replacing doer with defer (screen 16), involve only screen 23.

This program allows the user to define and initialize a file, enter data, query a file on any combination of fields, delete records and change field values of records.
FILE is the defining word for files. The PFA of a word defined by FILE contains the following parameters:

## byte

offset
0 starting block of file
2 maximum number of records for file
4 bytes/block
6 record length in bytes
8 current record number
10 address of list of fields for this file
FIELD is the defining word for fields. The PFA of a word defined by FIELD contains the following parameters:

## byte

offset
0 field width
2 byte offset from start of record
4 field type
The first record of each file (0 RECORD) is used for information regarding the length of the file (LASTREC) and the number of active records in the
file (\#ACTIVE). These items occupy the first four bytes of this record.

Screen 24 shows the file and field definitions for a sample application. Three parameters must be specified when defining a file: the starting block, maximum number of records and the record length. In the employees file definition, sixty-four is the record length, 100 is the maximum and thirty is the starting block.

When a field is defined, three parameters must be specified: field type, offset and field width. A width is specified even for numeric types for display formatting.

FIELDS compiles a list of CFAs of field words. The address of the start of this list is stored in field-list. The syntax for fields is: <file name> $n$ FIELDS. The file name executes and becomes the current file. The number of fields is then left on the stack to control the loop that compiles the list.

After the program has been loaded along with the sample file application (screen 24), type newfile employees. We are now ready to enter data into the employees file. Figure One shows a sample data entry session. Actually, the field entry prompts appear one at a time on the screen. When a field entry is terminated with a carriage return, the next field entry prompt will appear on the next line. After all the fields in FIELD-LIST have been entered, the user is asked if there is more data to be entered. The word nextrec in the ENTAY routine reclaims space occupied by deleted records.

Before querying a file, a display mode should be chosen by entering either the STEP or SELECT commands. STEP is the default mode. Figure Two shows a display using the STEP mode. It displays all the fields of each record found by the query. Records are displayed one at a time and the user is presented with the following options with each displayed record: modify a record, continue the query or quit. STEP does not require any arguments.

The select display mode allows the user to choose which fields will be displayed. This mode prints a heading with the names of the selected fields.

The field values for each record found are displayed under the corresponding field name in the heading. The syntax for SELECT is:
SELECT <file name>
<field2> $\ldots$ <field1>
The maximum number of selections (EXCERPTS) is five. This is arbitrary. More fields could be selected depending on the total number of characters of the selections. They should all fit on one display line. An example of a select display is given in Figure Three.

FIND is the end-user query word. It will search any combination of fields for each record in a file. The conditions are gr.than, ls.than, is and ISNT. The logical operators AND and OR are also used by the query. The maximum number of conditions ( $\mathbf{Q}$ \#) for the query is set for four. This could be increased, but since TiB will only accept eighty characters, I felt this was a reasonable maximum. A query requiring more than eighty characters could be input from a block using LOAD. If \#ARGS is modified to use a commandline delimiter, a number of query commands could be included on a block and loaded.
After the query command line has been entered, the program executes the next word in TIB which is a file name. It is now the current file. Next, the number of words that follow the file name in TIB is counted (\#ARGS). This number is incremented by one. If dividing this number by four leaves a zero remainder, the number of arguments is valid. The quotient is the number of conditions for this query. This value is left on the stack to be used by FOUND? and a-Arrays. Now the query arguments are stored. The search arguments (the values that are to be compared with the specified fields) are stored in targets. The maximum number of search arguments is thirty. Strings for numeric search arguments are converted by number before they are moved to targets by the word bring.

The file is searched, checking each record to see if it is active (not removed). If it is active, the query arguments are executed by FOUND?, which

processes the query arguments for each condition to see if the current record satisfies the conditions. After all the conditions have been tested, a flag is left on the stack. If it is true, then the query conditions have been matched by the current record and it will be displayed.

A word to list the entire file has not been included in this program. The entire file can be listed by entering a query with conditions that will be satisfied for all records (e.g., FIND EMPLOY.
ees name isnt xxx). The program includes very little error checking. If the user enters field names or conditions that have not been defined, the program aborts displaying the usual Forth system error message.

## Query Glossary

'OPEN Contains parameter field ad. dress of current file.
'FIELD Contains parameter field address of current field.

FIELD-LIST Address within parameter field of current file that contains address of list of fields for that file.

```
#%
GTEF क%
FTHT EMFlQYEES HOURS GF. THAN %O
NFIE " SWIFT
HOUFLY.WFOTE " 7.OO
HOUFE : さ%
ロEFT: ए%
FETUFN ta quit ESC to modify
eny fey to continue <space bar>
WF|E : WT1. कON
HOURLY--FATE : 7.7E
HIMFS: 4%
DEFT : DF
RETURN to quit Esए to modify
an. ter teremtimu% <ESC>
```



```
Fnter mame of fielo to be whanged
HOUFE
MOQRE % %
FETuF| to quit raE to imbuify
wy fey tr contame<space bar>
M&|E FIETH
```



```
HT|FO: 40
णएगT: &पए।
```



```
an* ley to wontwinue <CR>
#Mary awortect
```

Figure Two

LASTREC First byte of record number zero. It contains record number of last record in current file.
\#ACTIVE Third byte of record number zero. It contains the number of active records (not removed) in current file.

FILE File-defining word. When a word defined by FILE is executed, it places its parameter field address in 'OPEN.

FIELD Field-defining word. When a word defined by FIELD is executed, it places its parameter field address in 'rielo and leaves the address of the field on the stack.

FLD-wIDTH Contains the width of the current field. A field width is required for all field types. For numbers, the field width is required for display formatting.

FLD-TYPE Field types are 0, 2, 4 and 6, for text, single numbers, double numbers and dollar amounts.
table Defining word for execution tables of type-dependent functions. When executing, words defined by table use the current field width to select a function to be executed.


Figure Three
(ENTER) An execution table containing entry words for all field types. The words in this table all expect a field address on the stack at execution time.

DISPLAY An execution table containing display words for all field types. A field address is expected on the stack at execution time.

COMPARE An execution table containing words for comparing fields to search arguments. Words in this table expect two addresses on the stack and return $-1,0$ or 1 , for less than, equal or greater than.

DASHES Used for prompting input for record entry.

ENTER Prompts the user for a field entry. Accepts the input and stores the entry in the file.

REmOVED? True if record has been marked as deleted.
\#ARGS Counts the number of arguments remaining in TIB. Should be modified if block input is to be used for commands.

Q\# Maximum number of conditions searched for by query.
\#HITS Number of records found. In this application, \#HITS is only used as a flag, but it is easy to imagine other uses for it.

LOGICALS Array of logical operations (AND and OR) to be performed by query.

OPERANDS Array of field operands to be compared by query.

CONDITIONS Array of query conditions (GR.THAN, LS.THAN, IS Or ISNT).

TARGETS Address of start of storage area for search arguments.
+TARGET Uses index on stack to offset into targets.
bRING Execution table for words that bring the search arguments to TARGETS.
get-target Brings next word in tib to TARGETS using index on stack to offset.
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.HEADER A quick and dirty formatting word for the SELECT display mode.

SPREAD A quick and dirty formatting word which attempts to keep the displayed field values lined up under the field names in the display header.

SELECT End-user word for choosing fields to be displayed. Maximum number is five in this application. Can be changed, but the total number of characters of the fields selected should be less than eighty.

FIELDS Includes fields in the field list for a file after they have been defined. Expects the number of fields on the stack (<file name> n FIELDS).

STEP An end-user word to control display. All fields for a record will be displayed, one record at a time.

Newfile Initializes file by setting LASTREC and \#ACTIVE to zero.

Q-ARRAYS Uses the number on stack - which is the number of conditions
for a particular query - as a loop index to load query arrays with arguments from TIB. First entry in LOgicals is a no-op word.

FOUND? Compares fields with search arguments to determine if query conditions are satisfied.

FROM Executes the next word in TIB which is a file name.
(FIND) Examines every record in the current file, checking first to see if the

```
Screen # 8
I DOER/MAKE
: NOTHING :
: DOER CREATE ['] NOTHING \BODY , DOES\ & \R ;
VARIABLE MARKER
: (MAKE) R> DUP 2+ DUF 2+ SWAP E \BDDY ! % TDUP
            IF OR THEN:
: MAKE STATE I IF COMPILE (MAKE) HERE MARKER! 0,
        ELSE HERE [COMPILE]' `BODY! 1 STATE ! INTERPRET
        THEN : IMMEDIATE
: ;and CompILE EXIT HERE maRKER © ! ; Immediate
: UNDO ['] NOTHING >BODY [COMFILE] ' >BODY !;
Gcreen # 11
\ FILE words from FORTH Diaensions Vol. IV $5
VARIAELE 'OFEN I points to current file block
: REC& 'OPEN@ & + ; holds current record number
: LAYOUT I leave bytes/record-2, bytes/block-1
    COPEN E4 + 2¢ ;
: MAXRECS (-- n) 'OPEN e 2+ E:
: READ (nith rec, on stack, is made current)
    O MAX DUF MAXRECS < IF-NOT ." file error " QUIT THEN
    REC#!;
: RECORD (n -- a) \ leave address of n-th record
    LAYOUT 4/MOD 'OPEN + + ELOCK + ;
: ADDRESS (-- a) I leave address of current record
    REC# © RECDRD ;
: FIELD-LIST (-- a) 'OPEN & 10 + ;
: REC-LEN OPEN & b + ;
```


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```
Screen # 19
I QUERY mords elp 03sep85
: Q-ARRAYS ( n --) ['] NOTHING 0 LOgICALS!
    O DO I IF ! I LOGICALS! THEN
        DUP I OPERANDS! `BODY 'FIELD!
        I CONDITIONS ! I GET-TARGET LOOP;
: LOGIC ( i --) LOGICALS e EXECUTE ;
: OPERAND ( i --) OPERANDS & EXECUTE ;
: CONDITION ( i --) CONDITIONS E EXECUTE ;
: TARGET +TARGET ;
: FOlund? ( n -- f) 0 do I Operand I target I condition
        I LOGIC LOOP;
: FROM ' ( filename) EXECUTE ;
:(FIND) (n --) 0 #HITS ! LASTREC & 1+ 1
        DO I REMOUED? IF-NOT I REC# ! ( n) DUP FOUND?
            IF 1 IHITS +! CR .DISPLAY CR DELAY THEN THEN
        LOOP DROP #HITS & IF-NOT CR ." search failed" THEN;
Screen # 21
I FILE entry words elp 03sep85
: NEWFILE FROM O #ACTIVE! O LASTREC! ;
: FREE (-- rec\) LASTREC & 1+ I DO I REMOVED?
        If I leave then loof:
: NEXTREC (-- rec*) LASTREC & #ACTIVE &)
        IF FREE REC#! adDRESS REC-LEN BL FILL UPDATE
        ELSE LASTREC DUP E I+ DUP READ SWAP! UPDATE THEN;
: WRITE
        FIELD-LIST BEGIN DUF ?DUF WHILE EXECUTE ENTER 2+
        REPEAT DROP ;
: 3DOMN CR CR CR ;
: ENTGY #ARES 1`ABORT" needs filename" FROM
        bEgIN CLEARSCREEN 3DOWN NEXTREC WRITE I EACTIVE +!
        UPDATE BDOWN DONE? UNTIL
        SAVE-BUFFERS ;
screen # 23
I file display words elp 12sep85
    .EXCERPTS EXCERPTS BEGIN DUF & ?DUF WHILE EXECUTE
        .FIELD SPREAD 2+ REPEAT DROP ;
: SELECT \ u5age: SELECT <filename><fieldl> ...<tieldn>
        FROM EXCERPTS *ARGS DUP 5> ABORT" too many" 0
        dO ; DVER! 2+ LDOF O SWAP! MAKE DELAY NDOP {AND
        MAKE .DISPLAY .EXCERPTS ;AND MAKE HEADING .HEADEK;
    : FIELDS \ usage: filename n FIELDS fieldl field2 field3 ....
        HERE 5WAF 0 DO', LOOF 0, FIELD-LIST ! ;
: .MSSG CF ." fETLRN to quit ESC to modify" CR
            ." any key to continue" CR ;
: STEF MAKE DELAY .MSSG KEY DUP 27 = IF DROP MODIFY
        ELSE 13 = IF CK ." query aborted " ABORT THEN THEN
        {AND MAKE .DISFLAY . RECORD ;AND MAKE HEADINE NOTHING;
    STEF I default display mode
```


## Screen \# 20

1 QueRy words
Elp 03sep85
: FIND I end user query word
FROM
\#ARGS 1+ 4 /MOD SUAP ABORT" incorrect of arguments"
DUF Q\# > AEDRT" incorrect \# of arguents"
CR HEADING DUP Q-ARRAYS (FIND) ;
I usage: FIND EMPLOYEE DEPT IS PARTS AND HOURS GR. THAN 40
1 Other words
: DONE? ( -- t=no-more=entries)
CR ." any more? Y/N " KEY DUF EMIT ASCII $N=$;

Screen \# 22
1 Display header
VARIABLE EXCEFFTS 12 ALLOT | points to field to be displayed
: DASH-LINE CR 72 O DO ASCII - EMIT LOOP CR ;
: . HEADER EXCERPTS REGIN DUP © TDUF WHILE DUF B BODY
'fIELD! BODY; THAME DUF .NAME CE 31 AND FLD-WIDTH SHAF - ABS 1+ FFACES 2+ REPEAT DROF DASH-LINE CR ;
: SPREAD FLD-TYPE IF 'FIELD BODY XAME Ca 31 AND FLD-WIDTH SWAF - AES FLD-WIDTH + SPACES ELSE 2 SPACES THEN ;

Screen \# 24
| Application file and field definitions
6410030 FILE EMPLOYEES
ALPHA 120 FIELD NAME

* 21 6 FIELD HOURLY-RATE

SINGLE 252 FIELD HOURS
ALPHA 27 \& FIELD DEFT
EMplovees 4 fields name hourly-rate hours dept
record has been deleted. If it hasn't, it is checked to see if it matches the conditions specified in the query command line.

FIND End-user query word. Checks to see if an incorrect number of arguments has been entered in the query command line.
nextrec If the number of active records is less than LAStREC, the first deleted record (found by FREE) is used for the next entry. If there are no deleted records, the file is extended one more record.

Write Goes through the list of fields for the current file, prompting and accepting entries.

ENTRY A generic entry word for all files defined by file. The fields must be included in the field list (FIELDS).

EXCERPTS Address of start of the list of fields chosen by SELECT to be displayed.


## Across

1. A process control language
2. What a computer does
3. Pertaining to metal men
4. Changes in the flow of a program

## Down

2. What a bad program should do: abbr.
3. Liked by squirrels
4. Remove solid $\mathrm{H}_{2} \mathrm{O}$
5. Programmers in a frenzy
6. Type of transistor
7. Measure of resistance


Across
5. Exchange
6. Average
8. Make bigger
9. Indicates an address
11. Not odd
12. A computer noise

## Down

1. Creates machine code
2. What a programmer never wants to do
3. Give out
4. What stops the processor
5. Character $\qquad$
6. Look at memory location
7. Layer
(Answers on page 31.)

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# A Forth Standard? 

Glen B. Haydon<br>La Honda, California

What is a standard language? Natural languages evolve. Only after a word is used with a specific meaning for a period of time do dictionary editors consider including it. Many words have multiple meanings. Many definitions include examples of their use. Some words become obsolete or archaic. Languages are dynamic. They cannot be set in concrete. There is no such thing as "standard language." Dictionaries only record current usage.

Forth does not differ from any other language. It is evolving. That is the way Charles Moore designed it. He changed his kernel and application utilities almost daily. Many of you are aware that he includes a meta-compiler with most of his applications so he can easily recompile his kernel. It will be interesting to see what direction he takes now that he has cast his kernel in the Novix 4000 chip.

Before going any further, I would like to make a distinction between a kernel and a functional language. The Forth kernel is, in essence, the emulation of a hardware processor. The Novix 4000 is the implementation of a kernel in hardware. On the other hand, Forth as a functional language is built upon a kernel. It utilizes its extensibility to develop an operating system, compiler directives and utilities to solve problems. The functional language is a bridge between application requirements and the kernel. The beauty of Forth is the ease with which the necessary and sufficient functions can be added to a kernel.

The kernel usually includes between sixty and seventy hardware-related functions. There is little problem identifying these, but in actual hardware it has becomes obvious that some of the emulated functions are not optimal. Some of the problems were not anticipated by anyone.

The best example of a problem is the DO LOOP structure. The original figFORTH implementation requires a range in reverse order. What did the emulation do when a range crossed the
boundary of a signed number? Considerable error checking was added to the LOOP function in the 79-Standard definition. This proved to be a real boat anchor for speed nuts. This problem was addressed again in the 83-Standard and was improved. In the Novix 4000 the function was replaced by FOR NEXT. This function takes a count and decrements it to zero. The hardware requirements for speed dictated that a count-down register would work better and faster. Now the higher-level DO LOOP function becomes a part of the functional language, if it is going to be used. So the language changes.

With any Forth kernel, in hardware or emulated, it is an easy job to implement any desired dialect of functional Forth. Each vendor has his own idea of what should be included and what should be excluded. Each vendor provides a slightly different dialect of Forth. Most vendors make their kernel and the basic part of their functional Forth proprietary.

Let us review the public-domain versions of the primitive Forth functions. I started with the first public-domain version readily available - the figFORTH Model. The installation manual provided a verbal definition, and the several implementations clarified any possible misunderstandings. The system worked well. I did a moderate amount of programming with it.

Then came the 79-Standard. This was the result of about twenty Forth programmers who addressed some of the "problems" of the fig-FORTH Model. They did several things.

First, they changed the functional definitions for forty words previously defined in the fig-FORTH Model. Some of the changes were simply the use of an alias for the same function. Other changes were of a minor nature. The improvement to the compiler directive CREATE DOES $>$ was perhaps the most significant. The ability to write special compiler directives as part of an application program is unique to Forth among computer languages.

Second, the 79-Standard went beyond these functional changes. It in-
cluded a list of additional "Requirements" for any program adhering to the 79-Standard. In the Standard publication under Section 8, 'Use"':
"A Forth Standard program may reference only the definitions of the Required Word Set, and definitions which are subsequently defined in terms of these words ..."

This is patently ridiculous. At the November 1981 FORML Conference, I had an implementation of Forth which contained only the 148 words in the required word set. None of the members of the Standards Team who were there could do anything with the program. No vendor I know of has built a product in complete conformity with the restrictions imposed by the 79-Standard document.

About this same time, Robert L. Smith released and copyrighted a Forth-79 Standard Conversion. This publication consisted of a series of screens which could be loaded on a figFORTH Model. They would redefine the necessary forty words in the required word set. He admonishes the user to meet the other requirements of the 79-Standard.

Instead of conversion screens, I modified the compiler source code for the fig-FORTH Model to conform with the 79-Standard Required Word Set and made the additional functions required for a headerless operating system. This was a simple matter of changing a flag for the cross-compiler. I must acknowledge the efforts of Jerry Boutelle, who adapted his cross-compiler for the job and added many of the features. In a period of months two revisions were made. The resulting MVP-FORTH has remained stable for four years! The glossary All About Forth provides a reference to the common functions in public-domain implementations of FORTH up to that time.

Added to the MVP-FORTH kernel are a number of utilities and some supplemental definitions that will make this functional Forth almost completely compatible with Leo Brodie's Starting Forth. The differences are related to his use of a proprietary product (poly-

FORTH) which was supposed to be 79-Standard. Alan Winfield's The Complete Forth provides an excellent alternative tutorial.

Copyright protection of software is a continuing problem. The spirit of figFORTH was to put all of the source code and documentation in the public domain, asking only for appropriate acknowledgment. MVP-FORTH adopted the same spirit and placed all of the basic source code and documentation in the public domain. The contents of Volume 1 in the MVP-FORTH Series, All About Forth, are released without restrictions. Each entry includes a functional definition, indicates the source, an implementation, the usage in the MVP-FORTH kernel, an example with a note and a general comment. The general comment includes known differences in function among dialects.
As an interesting aside concerning the significance of copyrights, we had some correspondence with the publisher of Starting Forth. They claimed they had a copyright on all of the functional definitions included in their book. They claimed we could not include any of their functional definitions in All About Forth. I made an exhaustive study of prior functional definitions of the same words and was able to cite at least one prior definition for each word. Some of those prior definitions were also copyrighted and the publisher had failed to secure a proper release. So much for copyrights.

Other vendors approached the 79-Standard in various ways. Generally, their documentation has been excellent. I have always felt that the more implementations of Forth there are available, the more Forth will be used. By the time these products were on the market, the Standards Team was at it again and came out with the 83 -Standard. In my opinion, this was a great disservice to the advancement of Forth.

When the 83 -Standard was first available, I made a very careful comparison of the new functional definitions of the Required Word Set with those in the 79-Standard. The number of required words was reduced from 148 to 132. All but five had some
change in the functional definitions. No implementations were included as in the original fig-FORTH Model. In fact, some of the adopted functions had never been tested by the team.

In fairness to the members of the Standards Team, they are a dedicated group whose sole objective has been to improve and advance Forth. Many of the changes I found were simply attempts to clarify the wording of the previous standard.

However, they saw fit to change the functional definition of some words without changing the names. PICK and ROLL are examples. They required that the value on the stack be decreased by one from the value according to the 79-Standard. Thus:
: ROT 3 ROLL ; ( 79 -Standard)
: ROT 2 ROLL ; ( 83-Standard)
When you know of this incompatibility, it is easy to go through your code and change all the values to make it function. But I can see no improvement. Once a convention is adopted, stay with it.

I have no inclination to go through such a careful comparison again. Most of the changes made little difference. However, as has been observed by members of the Standards Team, most people don't do floored division. Forth has enough problems as it is. Why add to them with obscure changes? Forth needs stability.
In addition to the changes in the Required Word Set, similiar requirements to those cited above in the 79-Standard are included in the 1983 document. There is no way to verify the compliance of the many systems purporting now to be 83-Standard.
In the best spirit of Forth, Laxen and Perry have done an implementation of Forth which has become known as F83. It is unfortunate that this has been assumed to be the 83 -Standard. It goes far beyond the 83-Standard. It includes nearly 1200 words, and contains many excellent examples of problem solving with Forth. They provide full source code and shadow screens to assist the user. Unfortunately, there is

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no tutorial such as Starting Forth to go along with it. Every Forth programmer should be familiar with the many techniques these master Forth programmers have used.

Among the vendors, Laboratory Microsystems, Inc. has a version which is supposed to comply with the 83 -Standard. After finishing his implementation, Ray Duncan wrote a most interesting commentary on the 83-Standard which was published in Dr. Dobb's Journal. Other vendors have also implemented what they call 83-Standard Forth. Each of the vendors has excellent documentation for its particular implementation. A variety of other books on Forth are gradually appearing. Each is based on a specific Forth dialect, nuany of which are proprietary and copyrighted. However, many of the examples and ideas are portable to other Forth dialects with minimal effort. These books are a great help to the intermediate Forth programmer.

Already, some members of the Standards Team are soliciting suggestions for an 87-Standard. It is hoped that the FORML Conference this year will be able to address some of these recommendations.

I would humbly urge those interested in promoting the careful evolution of Forth to take a lesson from the pharmaceutical industry. Only after years in the chemical laboratory and more years of animal testing, are new drugs released for clinical trials. Only after all of the testing and trials have proven satisfactory are drugs finally released for general clinical use.

The Forth Modification Laboratory, FORML, is a fitting place for the laboratory development of modifications. The modifications should first be tried in the laboratory. Favorable results from such work should be submitted to clinical trial in the hands of vendors. Only by acceptance on the part of vendors should changes to a standard be adopted. But then it will not really be necessary: the modifications will have evolved into the common base of the functional Forth language. The standard will be established by common usage.

There is a recurring question of standard libraries. If people would publish their techniques, they could be adapted into most Forth dialects. But there is a reservation on the part of many authors. They want to have some return from all of their efforts. It is only reasonable that they be rewarded for their efforts.

Mountain View Press has found a partial answer to the problem. Namely, though some of their nine volumes are copyrighted, the contents are released for non-commercial use. At least the user can learn from the examples. It is highly likely that he will want to redo any algorithm in his dialect for his own application. Certainly it is not reasonable to let others reprint a book for profit as has been done with Volume 1 of the MVP-FORTH Series.

The current edition of Volume 3 in the MVP-FORTH Series is an example of the evolution of such thinking. The original text was written more than four years ago, and has been actively used since then. In 1985, author Phil Koopman agreed to a restricted copyright releasing it for non-commercial use. Each entry is modeled on $A l l$ About Forth and includes a functional definition, a high-level Forth implementation, an example with a note and a comment.

The local fig-FORTH community still objected: they could not use it because of the copyright, as open as it was. Some in the community have copyrighted their work and made no concessions to non-commercial use. This year, Phil Koopman released his work from copyright, with no restrictions. I hope more Forth authors will see fit to follow his example.

To argue about Forth standards is for those who have nothing better to do. Let Forth evolve like any natural language. Unlike other programming languages, it is easy to start over and meta-compile a new kernel. It is easy to build a new functional system.

Keep the FORML work active in the background. Encourage regional FORML workshops. As modern micro-
computers are becoming more powerful, something more than sixteen-bit address space is needed. How to incorporate this new hardware into the language presents several alternatives. None of the existing public-domain implementations address this problem. The existing standards are simply not mpatible with thirty-two-bit stacks.

Don't let the existing standards be an albatross to the language.

We have an urgent need for a stable language for beginners, for the management team entering new projects and for administrators new to the language. Let common usage provide a dynamic standard to meet the evolving needs. Let everyone participate.


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[^1](Continued from page 16.)
properly. STREAM-PROCESSOR: could also be used to implement the function of character translation by defining a character-parsing child. Other possibilities include a string search function for source screens. Implementing all these functions is made simpler and clearer through the added functionality afforded by a well-decomposed Forth kernel.

The dictionary look-up words shown in Listing One also make effective use of dual-CFA decomposition: the fail-ure-mode processing is factored into a child definition, which inherits a dictionary look-up function from the parent. So one word, the child definition, integrates and binds two related behaviors. While the child represents efficient factoring, the parent suggests a related family of words.

As shown in Listing One, the children of FAILING-LOOKUP: are ?COMPILEnUMBER; and ?INTERPRET-NUMBER; Both of these words represent incremental progress toward their parent functions, cOMPILE-WORD and INTER-PRET-WORD. Note also that these stringhandling functions need not be expanded any further to produce a workable system (as will be shown). To expand them any further would produce undesirable crossover into the domains of other families of words.

As defined in Listing One, streamPROCESSOR: actually combines three behaviors into each of its children. The parsing loop is inherited by the children, but it also contains a vectored execution that specifies the processing after each word is parsed. The child merely specifies the version of WORD to be used within the shared word-parsing loop. (See tib-process, blk-process, tIB-WORD and block-word.)
The flexibility needed to switch from compiling a word to interpreting a word at run time (and vice versa) requires the use of a vector. The left and right bracket definitions must reinitialize the vector. Since the brackets may occur amidst an input stream, the action of the children of STREAM-PROCESSOR: is also variable midstream. To expand Listing One to include bracket definitions, you could use:

```
: l (--)
    192 STATE!
    , COMPILE-WORD
    CFA PROCESS-WORD' 1 ;
: [ (--)
    O STATE!
    , INTERPRET-WORD
    CFA PROCESS-WORD' 1 ; IMMEDIATE
    Finally, the Forth functions normal-
ly performed by quit and interPret
can be easily constructed as a single
definition:
: INTERPRET
    [COMPILE] [
    RP! BEGIN
    CR QUERY TIB-PROCESS
    STATE @ 0= IF
        ."OK" THEN
    AGAIN ;
```


## Early Impressions

The relative newness of dual-CFA decomposition has not prevented me from forming opinions regarding its most suitable use.

I have some reservations about the implementation of deferred definitions (DEFER:). I prefer to see a closer relationship between the two functions bound together through dual-CFA decomposition. In DEFER:, the parent definition provides a compiler-extending behavior and the child definition forward references to an arbitrary function.

I favor failing-number: and FAILING-LOOKUP: as examples of how dual-CFA decomposition techniques should be applied. I appreciate how closely united the parent and child definitions are: the parent look-up function is made more specific by the failure mode processing provided by the child. In actual use, the child refers to both functions as if they were a single, undecomposed function. Yet because they are decomposed, you are free to define new children without restating the parent function.

The demystification of Forth would be a welcome by-product of a more clearly and more fully decomposed kernel, if one should ever find its way into widespread use. Some evidence of this can already be seen in Listing One: (1) The end-of-input-stream detection function is within the parent stream-
processing function, not hidden in a definition of NULL. (2) The STATE variable is less central to one's comprehension of Forth - the interpret and compile functions are explicitly separate, even though they still share a common word-parsing loop. (3) Words that manipulate input streams are more easily distinguished from words that perform interpreting or compiling actions.

On the other hand, programming became more difficult than before. The program code in Listing One required subtle but definite changes in my programming style. Many times, I had to abandon a particular approach in search of something more intuitively obvious. However, the development process did fine tune my perception of the problem along functional lines.

The functional areas of concern required clearer identification at the outset. Next, each of these functional areas had to be well decomposed. Finally, refinements were made so that the stack effects of all functionally related subsets of words belonging to a particular family remained consistent. The comment header shown in Listing One also helped.

Throughout development, a continual effort was necessary to prevent subsets of words from wandering into the domain of another family of words. I cannot overemphasize the point that this kind of programming demands a clearer delineation of definitions along functional lines. Hybrid words must be acknowledged before useful dual-CFA decompositions can be found (such as the effort surrounding WORD).

## Conclusions

The examples shown of dual-CFA decomposition have helped illustrate some of the advantages possible with this methodology (see iterator: ${ }^{1}$ as well). A summary of the advantages includes:
(1) Better organized definitions, particularly along functional lines, in-
creasing the ease with which Forth source code can be read and understood.
(2) Increased emphasis on more complete decomposition, resulting in a richer programming environment and increased productivity.
(3) Decreased likelihood of programming error and system crashes, through elimination of many environmentally dependent behaviors.
(4) Decreased need for passing flag parameters on the stack, as well as a corresponding decrease in the number of conditional-behavior words (control-flow constructs such as IF THEN now are factorable and need appear only once per function - even if the function is decomposed).
(5) Increased memory compactness for compiled applications.

Also, modern innovations associated with new programming languages or operating systems may be more easily implemented. Examples might include object-oriented modules, relocatable modules and "piping" capabilities for stream-processing modules. These areas are generating more and more interest lately. Dual-CFA decompositions can bring each of these areas of programming interest within closer reach.

Someday, perhaps, the Forth dictionary will be mostly a library of forms ${ }^{2}$ or general algorithms, from which a programmer compiles more specific instances of each algorithm to accomplish a particular task. If this happens, each issue of Forth Dimensions may include many practical applications. Each would be derived easily using provisions already included in the Forth dictionary.

## References

1. Elola, Mike. 'Dual-CFA Definitions,'' part one, Forth Dimensions VIII/2.
2. Luoto, Kurt. "Procedural Arguments," Forth Dimensions VI/2.

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# Windows for the TI 99/4A 

Blair MacDermid Fort Wayne, Indiana

The Forth routines described here provide plotting of algebraic functions in a choice of five windows located in different positions of the display screen. These windows can be displayed simultaneously. Four of the windows can be located in the four quadrants of the display screen. The fifth window occupies most of the full screen. The program computes the coordinates of the plotted function, appropriately scaled to fit within the selected window.

These routines are a by-product of the group project undertaken as a learning exercise by members of the Fort Wayne FIG Chapter. The primary objective of the exercise was to allow the group members to participate in development of a useful Forth program, with efforts partitioned among members according to their skills. The simpler elements of the program were assigned to beginners. We also hoped to learn how well Forth would serve in a multiple-programmer task.

Ed Harmon, the chapter's guru, selected the ACM SIGGRAPH CORE Standard as a useful tool. He provided us with a model written in the UCSD p-System Pascal (see the Journal of Pascal, Ada, Modula-2, May/June 1984, page 19). Ron Bishop, president of the local TI 99/4A Users Group, completed the program using thirty screens and integer arithmetic. That implementation provides the freedom to locate and define the size of a number of viewports (i.e., windows) to be displayed simultaneously. The size and location are continuously adjustable.

Here I have defined a simpler version of the program, using only six screens. It does not provide the degree of freedom intended to be part of the ACM standard. However, it provides a useful choice of window locations and sizes. The program exploits the excellent graphics capabilities of the TI 99/4A using the SPLIT2 mode and the TI-FORTH words DOT and LINE, as well as the TI 99/4A's floating-point routines (which I used to plot functions
that contain the transcendental functions, e.g., sine, tangent, logarithm).
It will be useful to refer to the Forth screens $30-35$ in the discussion that follows. These screens provide a useful utility but can readily be modified and expanded to include different elements of the Pascal model of the full ACM SIGGRAPH CORE Standard.

## Screen 30

Lines 1-9 define the required variables. The variables associated with the horizontal axis use X , as is common practice. XMIN and XMAX represent the minimum and maximum values of the real-world function to be plotted. VL and VR represent the left- and righthand viewport coordinates in pixels, referenced to the TI screen display. Similarly, vвот and vTOP represent the bottom and top coordinates of the viewport.
The function $\mathrm{Y}=\mathrm{f}(\mathrm{X})$ is computed using X as the independent variable in the world coordinates. $\mathbf{x D}$ is the corresponding variable referenced to the display screen coordinates. Similarly, YDB represents the display screen coordinate corresponding to Y of the world coordinates.
The variable YDB warrants further explanation. The TI 99/4A screen display uses coordinates that reference the upper left-hand corner of the screen as the 0,0 point. I found this confusing, since it is normal to use the lower lefthand corner as the origin when plotting functions. So I invented yDB to allow me to readily handle the necessary mathematics. Subsequently, when specifying the coordinates to be plotted on the screen, I use the constant YTio equal to 191 (see screen 31 , line 2 and screen 35 , line 5) to make the necessary corrections for the TI 99/4A screen coordinates. This is justified by the following relation:
$($ TI's Y) + YDB $=$ YTIO $=191$
Therefore,
(TI'S Y) $=$ YTIO - YDB
The variables $\mathbf{K X}$ and $\mathbf{K Y}$ are scaling factors modifying the world coordinate
variables to fit the selected viewport dimensions.

The variable ydbara is an array to store 200 computed values of YDB. The TI 99/4A shares some of its display facilities with the floating-point routines; to avoid any difficulty in this regard, I chose to compute the values of the plotted function (see screen 34) before using the screen to display the function (see screen 35).

All of the variables discussed above represent integer values. This assumes that the selected minimum and maximum world coordinate variables will be integer values, the normal thing to do. Certain of these variables will require floating-point representatives in the computation routine of screen 34. These variables have been prefixed with an F as in $\mathbf{~} \mathbf{X X}$, $\mathbf{F K X}$ and $\mathbf{F K Y}$.
The words kxcalc and kycalc specify the computation of the scaling factors $\mathbf{K X}$ and $\mathbf{K Y}$, as well as the floating-point equivalents FKX and FKY. It may be helpful to display the mathematical definitions of these variables:

## $K X=(V R-V L) /(X M A X-X M I N)$ <br> $K Y=($ VTOP-VBOT) $/(\mathrm{YMAX}-\mathrm{YMIN})$

TI-FORTH uses the word s->F to convert integer values to floating point.

## Screen 31

This screen contains the definition of the viewport size and location. The word PICKPORT selects the viewport corresponding to the number (one through five) placed on the stack before executing the word. The numbers one through four select a viewport, size $100 \times 50$ pixels high, at locations in any of the four quadrants of the display screen. A value of five selects the largest viewport, 200x100, occupying most of the available screen display. There is sufficient space to allow a height of 190 pixels; however, it is easier to interpolate values of Y with the height of 100 pixels. Of course, other viewport dimensions and locations can readily be specified by substituting different numbers. (I find

the word PP useful in exercising the program, since it relieves me of the need to type PICKPORT, whose length is dictated by the desire to write readable code.)

## Screen 32

The words Ylines and xlines use TI-FORTH's LINE to draw vertical and horizontal lines at useful increments. The word GRID uses these words to superimpose on the viewport a grid to expedite interpolation of values of the displayed function. The words $\mathbf{x o}$ and yo are abbreviated versions of the words locating 0 axes for both X and Y .

## Screen 33

The words C1, C2 and C3 specify different parameters for the world coordinates of the function to be plotted. The word function specifies the function to be plotted, in this case $\sin x / x$, defining the Fourier spectrum of the rectangular pulse waveform. Sufficient space is available on this screen to substitute another definition of the word function. Notice, however, that the definition must use floating-point representation.

## Screen 34

The word ARAYDB specifies the computation of the values stored by YDB in the array ydbara. It also causes the word wORKING!!! to be displayed on the screen so that the user will not assume his computer has contracted amnesia while executing the calculations. Note that lines $2-5$ perform calculations in integer arithmetic, and the results are converted to floating point by line 6. Line 8 contains the word FUNCTION, and the resulting computation is con-
verted to integer by the word $\mathbf{F - > S}$ (a single-precision integer value in two bytes). The word CY is my convenient macro for ARAYDB.

The DO LOOP increments the current value of XD by one pixel from VL to VR. Lines 3-4 compute the corresponding value of the world coordinates. This value is converted to floating point and is divided by the floating-point representative of the scaling factor $\operatorname{FKX}$. The result is placed on the stack, and a copy is stored in FX where it can be used in more complicated functions requiring different powers and functions of $\mathbf{X}$.

The computed value of FUNCTION is multiplied by the scaling factor $\operatorname{FKY}$ to define the corresponding value YDB for the display screen. The resulting computation for each increment in $\mathbf{X D}$ produced by the DO LOOP is stored in the array ydBara.

## Screen 35

The word PLOTY uses another loop to increment XD in one-pixel increments and selects the appropriate element of the array ydBaRa to plot the function on the screen using the TI-FORTH word DOT. (PY is my macro for initiating the plotting routine.)

## Final Notes

The definitions used assume the SPLIT2 graphics mode if the TI 99/4A is used. TI-FORTH is a fig-FORTH extension, but the words DOT and LINE are probably machine dependent. However, it is reasonable to assume the screens could be modified to work on a different Forth implementation. Both the Apple II and the IBM-PC have graphics capabilities providing pixel resolution.

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# Getting Started with F83 



Greg McCall<br>Werrington, NSW, Australia

The documentation with F83 is in F83.COM and in the shadow screens that are part of the source files that come with F83. At first glance, the thought of sifting through hundreds of kilobytes of shadow screens is bewildering, to say the least. Just to get you started, I have put together a summary of how to use the file words and how to edit these files. This relates to the CP/M-80 version of F83, but as far as I know it should be similar to other versions of Laxen and Perry's F83.

This Forth can have two files open at once. One file is called the current file. This is the file used by all normal reads and writes. You would normally edit or load from the CURRENT file. The other file is called the From file. This is a second file you may have open for reading only. For example, if you currently are working on a file (i.e., loading and editing), and you wish to load some screens from another file, then you may open a FROM file and load screens from it without changing the CURRENT file. Following is a description of some useful file words:

CREATE-FILE (S n -- )
Creates a new file containing $n$ blocks.

## 10 CREATE-FILE TEST.bLK

opens a file called test.blk and writes ten blank screens to this file. The file is then closed.

FILE? Prints the name of the CURRENT file.

DIR Prints the directory of the current drive.

Open the following file name and make it the current file, e.g., OPEN TEST.BLK

Make the next word in the input stream the FROM file and OPEN it. It then sets the current vocabulary to FILES.

LOAD

CA

COPY
In the FORTH vocabulary, copies a screen in the CURRENT file. In the files vocabulary, copies a screen from the FROM file to the CURRENT file. In the SHADOW vocabulary, copies a screen and its shadow in the Current file.

## CONVEY (S from to -- )

In the fORTh vocabulary, copies a set of screens in the CURRENT file. In the FILES vocabulary, copies a set of screens from the FROM file to the CURRENT file. In the SHADOW vocabulary, copies a set of screens and their shadows in the CURRENT file.

## HOPPED <br> HOPED

A variable containing the number of screens to skip when copying with CONVEY.

U/D

то
A variable containing the direction of the screen move using CONVEY. +1 is a forward screen move and -1 is a backward screen move.

Sets up the variables HOPPED and U/D. Used as first-source last-source то first-destination CONVEY
The F83 editor uses the same words as the editor in Starting Forth by Leo Brodie, with some additions such as the word NEW which allows replacement of

In the FORTH vocabulary, LOAD will load screens from the current file. In the fles vocabulary, LOAD will load screens from the FROM file. So while we have a file as the current file, we can still open another file by making this second file the FROM file and loading from it, e.g., FROM TEST.BLK 10 LOAD

Copy a screen to its shadow.
(S from to -- )

multiple lines. To get the editor going correctly, you should look at screens 28 30 and 88 of UTILITY.BLK which hold the terminal-dependent routines. You can select your terminal - or see if any of the routines are the same as those of your terminal - or write your own routines. The terminal words patch the words AT, DARK, BLOT and -LINE to suit your terminal. While we are looking at patching the editor, you could remove the backslash in line 14 of screen 24 so that (WHERE) is patched into WHERE and, if you have a real-time clock, then you could change GET-ID in screen 23 so as to have the ID supplied when the editor is first invoked. These screen numbers refer to the $\mathrm{CP} / \mathrm{M}-80$ version of F83. To find where the source screens are for the editor in your Forth, type VIEW AT which should give you the second source screen of your editor. Now just look through the editor's screens for the required words.

A summary of the editor commands follows:

TOP Go to the top of the screen.

C (S n -- )
Move n characters, right or left.
T (S n -- )
Go to beginning of line $n$.

| BUFS | Displays the contents of <br> the insert and find <br> buffers. |
| :--- | :--- |
| KEEP | Places the current line in <br> the insert buffer. |
| K | Exchanges the contents <br> of the insert and find <br> buffers. |
| W | Write all changes to disk. |
| N | Move to next screen. |
| B | Move back a screen. |
| Alternate between a |  |
| screen and its shadow. |  |

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<text> Represents the text following the command. If <text> is just a carriage return, the contents of the insert buffer is used in place of the missing <text>.

| $\mathbf{I}$ <text> | Inserts <text> on the <br> current line at the cursor. |
| :--- | :--- |
| $\mathbf{O}$ <text> | Overwrites <text> onto <br> the current line. |
| $\mathbf{P}<$ text $>$ | Replaces the current line <br> with <text> and blanks. |
| $\mathbf{U}<$ text> | Inserts a line under the <br> current line. |
| $\mathbf{F}<$ text> | Finds the <text> and <br> leaves the cursor just past |
| it. |  |

Searches for <text> through all screens from the current one up to screen $n$. Each time a match is found, n remains on the stack until screen $n$ is reached. To continue the search, just type S until screen $n$ is reached.
$\mathbf{R}<$ text $>\quad$ Replaces the text just found with <text>.
D <text> Finds and deletes the text.
TILL <text> Deletes all text on the line from the cursor up to and including $<$ text> .

JUST <text> Deletes up to, but not including, <text>.
KT <text> Puts all text between the cursor and <text> inclusive into the insert buffer ('keep-till'').

E

X
SPLIT

JOIN

WIPE

G
Gets a line from another screen and
inserts it in front of the current line. In the shadow vocabulary, G gets a line and its shadow. In the files vocabulary, $\mathbf{G}$ gets the line from the $\mathbf{F R O M}$ file.

## BRING

(S screen first last -- )
Brings several lines from another screen and inserts them in front of the current line. In the shadow vocabulary, BRING gets a range of lines and their shadows. In the FiLes vocabulary, bring gets the lines from the from file.
NEW
( $\mathrm{S} \mathrm{n}-\mathrm{-}$ )

Moves the terminal's cursor to the start of line n and overwrites until the line has a null input, i.e., just a carriage return.

QUIT

DONE

ED
Re-enters the editor. It clears and reinitializes the display, and begins automatic redisplay of the screen.

EDIT ( $\mathrm{S} \mathrm{n}-\mathrm{-}$ )
Sets SCR to n , then uses ED to start editing.

This should enable you to copy and edit screens with Laxen and Perry's F83. The best way to get the entire documentation on this Forth is by printing out all the source files. If your printer can print at least 132 characters per line, then look in your printer's manual for the characters needed to put your printer in this mode. My FAX-80 needs a control-O to set the condensed mode. I define a word FAX-80 to send this code and then patch it into the DEFERed word INIT-PR, i.e.:
: FAX-80 CONTROLO EMIT ; , FAX-80 IS INIT-PR

To print the entire file, you use the word LISting. For example, to list META80.BLK, I would type:
OPEN META80.BLK LISTING

# Batcher's Sort 

John Konopka Mitaka Shi, Japan

Quicksort is often suggested as a sorting algorithm because of its speed. The reputation for speed is well deserved but Quicksort has other features which may make it difficult to use. An alternative sorting method discovered by K.E. Batcher in $1964^{1,2}$ is a little slower than Quicksort but is more robust and avoids most of Quicksort's pitfalls.

One problem with Quicksort is its variable performance. It is usually stated that Quicksort requires about $\mathrm{N} \log \mathrm{N}$ operations to sort N items (base 2 logarithm). This is an average result which depends on the input data being random. In other cases where the data is already ordered in some way, then Quicksort may require as many as $\mathrm{N}^{2}$ operations to sort N items. This is as slow as a Bubble sort. Thus you don't know from one execution to the next just how long a sort will take. Extra code can be added - complicating the algorithm - to handle some, but not all, of the time-consuming cases. Quicksort also varies in its use of space. Every branch in Quicksort creates one stack entry (the number of words per stack entry is implementation specific) on the return stack (if, as usual, recursion is used). Normally, a maximum of about $\log \mathrm{N}$ stack entries are created. However, in degenerate cases this number may approach N . When sorting illordered data you may find your program running out of room with unanticipated consequences.

A second source of trouble with Quicksort is that it is difficult to implement. Quicksort is generally presented in a recursive form. If recursion is not available you must implement this yourself. You can, at the expense of more complicated code, implement a non-recursive version ${ }^{3}$. To limit, but not eliminate, the number of cases requiring much time or much stack space more code can be added, again increasing the complexity of the algorithm. The final implementation problem is how to test it. Because the operation of the algorithm is data

## Sample portions of link map data.

| OVLY | SEG | SIZE |  |
| :---: | :---: | :---: | :---: |
| 1 | 1 | 1833 | \#************************************* |
| 1 | 2 | 1130 | ************\#\#********* |
| 1 | 3 | 1972 | \#**************************************** |
| 1 | 4 | 2245 |  |
| 1 | 5 | 1696 | *********************************** |
| 1 | 6 | 2495 | **************************************************** |
| 1 | 7 | 2402 |  |
| 2 | 10 | 1499 | ******************************* |

## Before sorting

| QULY | SEG | SIZE |  |
| :---: | :---: | :---: | :---: |
| 1 | 21 | 2621 |  |
| 1 | 27 | 2618 |  |
| 1 | 25 | 2509 | \#*************************************************** |
| 1 | 6 | 2495 |  |
| 1 | 13 | 2485 | **************************************************** |
| 1 | 17 | 2469 | ****** ${ }^{*}+* * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * *$ |
| 1 | 14 | 2443 |  |
| 1 | 7 | 3402 |  |

After sorting
Figure One

| 100 random numbers before and after so |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 14820 | -10904 | 29081 | -30212 | 3226 | 16975 | -6865 | -31694 | 28585 | 29040 |
| -22503 | 25399 | 24896 | -27251 | 228 | 29720 | -10403 | 11702 | 84 | -13 |
| 1293 | -17882 | 11160 |  | -28685 | 21788 | 364 |  |  |  |
| -24843 | 27148 | -12 | 17090 | 28362 | -16741 | 249 | 12214 | 32405 | 78 |
| 31832 | -7663 |  |  | -7458 | -12676 | -7101 |  | -6936 | 360 |
| -15913 | -13499 | -27433 | 14612 | -8610 | -26152 | -9637 | -19365 | 6962 | 143 |
| -31048 | -19079 | 711 | 13083 | -16616 | -148 | 15938 | -19628 | 19793 | 20656 |
| 22997 | 32032 | 18638 | -9148 | 1954 | 8 | 1 | -17276 | 11578 | -16357 |
| 17601 |  | -36 | -205 | -14952 | -5764 | 26437 | -28174 | -474 | -22 |
| -32679 | 6053 |  | -1349 |  |  | 26 |  |  |  |
|  |  |  |  |  |  |  |  |  |  |
| 28362 | 271 | 264 | 25399 |  | 229 | 218 |  |  |  |
| 17601 | 16792 | 159 |  | 14612 | 13083 | 12214 |  | 115 |  |
|  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  | 364 | 249 | 74 | -1267 | -1349 |
| -3600 | -3684 | -5764 | -6865 | -6936 | -7101 | -7458 | -7663 | -8360 | -8610 |
|  |  |  |  | -12676 | -13499 | -13959 | -14024 | -14444 | -14840 |
| -14952 | -15913 | -16357 | -16616 | -16741 | -16975 | -17090 |  | -17882 | -19 |
| -19365 | -19628 | -19793 | -20587 | -22250 | -22334 | -22503 | -24843 | -26152 | -263 |
| 27251 | 27433 | -28174 | -28685 | -30212 | -31048 | -31694 | -321 | -32405 | -326 |

Figure Two
dependent you may have sleeping bugs which only awaken when presented with rightly ordered data. See the Sedgewick and Knuth references for more information about Quicksort.

Batcher's sort suffers none of these problems. It iterates the same way every time, calculating the same pairs of indices regardless of the data presented for sorting. It sorts in place, requiring no buffer space, and it places no unusual demands on either the return or data stacks. Furthermore, it is easy to implement, requiring only one screen of Forth code. Recursion is
not required. Finally, because it is simpler there are fewer things to go wrong. It is thus easier to test and easier to trust. Once you have it working for one set of data it is likely to work well afterwards.
The cost for this robustness is time. Quicksort requires, on average, about NlogN operations. Batcher's sort requires less than $(\mathrm{N} / 4) \log \mathrm{N}[(\log \mathrm{N})+1]$ iterations. The difference is less than $(\log N+1) / 4$. As an example of what this means in terms of normal array sizes Quicksort should be, on average, about two times faster when sorting 1024 items.

This does not take into account any time difference for one iteration between Quicksort and Batcher's sort. The clincher is the phrase "on average." Depending on the input data, in some cases Batcher's sort may in fact be quicker than Quicksort. In any event The absolute difference in time will probably not be large. For example, using no code words I can sort 512 names on a DEC LSI $11 / 23$ in twelve seconds. In this case the cost for using Batcher's sort is certainly tolerable.

Batcher's sort has one more interesting feature which someday may let it far outpace Quicksort or any other sorting method, in terms of speed. Looking at the code you can see three nested loops. At every iteration of the innermost loop INNER-LOOP the pairs of keys which are compared are completely independent. Thus a parallel computer could implement the inner loop in one step for really fast sorting. The number of iterations in this case is just $(1 / 2) \log N[(\log N)+1]$. This is just fifty-five iterations when processing an array of 1024 items.

## Implementation

The Forth code for the sort is displayed in screen 2 . While the code is not particularly complex, the operation of the algorithm is not obvious. See Knuth for further details. The program uses seven constants: TT, PP, DD, NN, $\mathbf{R R}, \mathbf{Q Q}$ and KC. These names were chosen to be consistent with the description of the algorithm given by Knuth. QQ can easily be carried only on the stack but I made it explicit for easier reading. Constants are used rather than variables, as the data is accessed much more often than it is set. $\pi$ stores a parameter which determines the sizes of the outer loops. It is calculated in seLect-t. PP drives the outermost loop, QQ drives the next nested loop. These loops are driven by dividing the loop counter by two rather than by incrementation as in DO LOOP. RR, NN and DD are used to calculate indices to keys. When sorting N items, this routine generates indices in the range from zero to $\mathrm{N}-1$. The actual output of the program is this sequence of number pairs. Implementation-

```
Screen #2
    \ BSORT K. E. Batcher's sort. From Knuth, vol 3.
    0 CONSTANT TT 0 CONSTANT RR 0 CONSTANT DD 0 CONSTANT PP
    O CONSTANT NN 0 CONSTANT QQ 0 CONSTANT KC
    : KEY COMPARE KC EXECUTE ;
    : SELECT-T NN 15 0 DO DUP I 2**N <= IF DROP I LEAVE THEN LOOP
                        1-14 MIN ' TT ! ;
    : INNER-LOOP NN DD - 0 DO I PP AND RR =
                        IF I DUP DD + KEY_COMPARE THEN LOOP ;
    : Q-TEST QQ PP <> IF QQ PP --, DD ! QQ 2/ 'QQ !
            PP ' RR ! O THEN ;
    : QRD-SET TT 2**N ' QQ ! 0 ' RR ! PP ' DD ! ;
    \ n-- n is number of items to sort. n must be positive.
    : BSORT ' NN ! SELECT-T TT 2**N ' PP !
        BEGIN QRD-SET QQ
                                    BEGIN INNER-LOOP Q-TEST UNTIL
                                    PP 2/ DUP ' PP ! 0= UNTIL ;
Screen #6
    BSORT example. Sort array of integers.
    0 CONSTANT X1 0 CONSTANT X2 CREATE DATA 200 ALLOT
            --- Load array DATA with random numbers.
    : INIT-DATA 100 0 DO RANDOM DROP I 2* DATA + ! LOOP ;
    \ --- Exchange entries pointed to by X1 and X2.
    : SWAP-DATA X1 DATA + @ X2 DATA + @ X1 DATA + ! X2 DATA + ! ;
    N M --- Compare and maybe exchange Nth and Mth entries.
    : COMPARE-AND-SWAP 2* ' X1 ! 2* ' X2 ! \ Save pointers
        X1 DATA + @ X2 DATA + @ > X2 \ Compare values
        IF SWAP-DATA THEN ; \ Exchange if misordered
    \-- . R is defined in 79-Standard Reference Word Set.
    : LIST-DATA 100 0 DO I 2* DATA + @ 7.RI 1+ 10 MOD 0=
                IF CR THEN LOOP;
```

(Screens continued from page 36.)
SCR \#34
\#R ." SCR\#S4 " ( ARRAY YDB CALC FLTE POINT BWM 7/7/85 )
: ARAYDB CLS." WORKING !!! " ( --- YDB[iJ)
VR VLGDOIXD !
XMIN VL KX


FDUP FX F! (stores current $f x$ )
FUNCTION FKY FE F* ( $f x$ - - fy*fky)
F->S (fky*fy --- ky*y)
YMIN K KY © - VBOTE
YDBAFA I VL $\underline{a}-2 *+$ ! (store ydb in array ydbara)
LOOP ;
( COMFUTES VIEWPORT FEPRESENTATIVE OF WORLD $Y$ )
4
5 : CY ARAYDG: -->
SCR \#35
CR ." SCR\#35 " ( EWM PLOTTING UTILITES 7/7/85 )
: floty

I VL (a-2*
YDBARA +
YTID SWAP - XD G SWAP DOT ( FLOTS NEXT PT,
LOOF ;
: FY CLS FLoty;
10
11
12
13
14
specific code uses these pairs of numbers to point to the items to be sorted, then does the compare and possible exchange. If, for example, you were sorting a list of names and the output was 1 and 5 , then your implementationspecific word would compare the first and fifth names in the list and exchange their positions if they were misordered. The execution address of this code is stored in constant kc. The word KEY_COMPARE accesses that constant and executes the word whose address is stored there. By this vectoring, the sort routine is separated from the data being sorted so you can use the same sort routine for all applications. To use the routine, put the execution address of your compare code in constant KC, put $\mathbf{N}$ (the number of items to be sorted) on the stack, then invoke BSORT.

## Application Examples

Screen 6 shows an application which sorts an array of random data. The array is initialized with a random number generator ${ }^{5}$ by invoking INIT-DATA. If a random number generator of some kind is not available you can load the array with an editor, using, (comma) to enter integers picked from your imagination. LIST-DATA will type the data on a terminal. To sort the data put the execution address of COMPARE-ANDSWAP in constant KC, then put 100 on the stack and invoke bSORT:

## FIND COMPARE-AND-SWAP ' KC ! 100 BSORT

Now you can use LIST-DATA to see the effect of sorting. This simple example is useful for verifying the operation of

## BSORT.

As another example application, I use this routine to sort the vocabulary names in the Forth dictionary. Code specific to my system first scans the dictionary and builds an array of addresses. Each entry points to the name field of a Forth word. The length of the array is the number of words in the
dictionary. The comparison word deposited in KC takes two indices from BSORT and using these pointers compares two names in the dictionary, then exchanges the addresses stored in the array if the names are not in alphabetical order. This comparison word must not only know how to compare strings alphabetically but it must be able to strip out special bits such as the immediate flag, and it must be able to determine the length of the name. After sorting, I write the names to a text file and then use an editor to make glossaries for documenting applications. See the paper by Baden ${ }^{4}$ for another example of sorting vocabulary names in the Forth dictionary.

In another case I use this routine to sort information about a large Fortran program. When the program is compiled and linked, a map is generated giving, among other information, the size of each of the program overlays. The size of the program in memory is determined by the largest segments; thus, to reduce the memory requirements one needs to know which are the largest segments and how they differ from the second or third largest segments. I wrote one routine to scan the map and extract the size information. For sorting, the word deposited in KC compares these sizes numerically and exchanges them if they were out of order. Figure One graphically shows the results before and after sorting.

In the near future, I have two more sorting applications in mind. One is in an application I wrote called "Card File." This is a software version of a box of $3 \times 5$ cards. In this case I will first create an array of pointers in memory indicating which cards I want to list on the printer. Then I will sort this list using BSORT. The most natural order would be to alphabetize the cards according to the first word on a given row of the card. The application-specific word which would be deposited in KC would have to know how to extract this information from the cards, then do the compare and swap pointers if needed. The second application I have in mind is in x-ray spectroscopy. I now have Forth words which create directories of file names of stored x-ray
spectra. It would be helpful to sort these directories in various ways. Just by changing the compare word deposited in constant KC I will be able to sort the directory according to file name, date, number of elements in the spectra or even according to the atomic numbers of the elements which generated the spectra.

From these few examples you can see that almost anything can be sorted. All you need is a word which knows how to compare two items in a list and exchange them if they are misordered. If the items are small and easy to move, then you can exchange the positions of the items themselves. If it is costly to move the items, as in the case of diskbased data, it is better to keep a list of pointers and just exchange the pointers.

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5. Doyle, William T., "A Portable Forth Random Number Generator', Journal of Forth Applications and Research, vol. 1, no. 2, 1983.

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