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FEATURES

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What Forth Can Learn From C

A recent ad for an Amiga version of Multi-FORTH from Creative Solutions reminded me of Forth. I had almost forgotten about Forth because recent issues of Computer Language and Dr. Dobb's Journal of Software Tools treat Forth as if it no longer exists.

Wednesday night I went to a well-attended Amiga owners meeting to see a demonstration of a brilliant new piece of Amiga software from a company called B.E.S.T. (P.O. Box 852, McMinnville, Oregon 97128). The software will allow small businesses to do all their bookkeeping using an Amiga. B.E.S.T. began last fall using Lattice C and recently switched to Manx's Aztec C. When I asked these brilliant programmers if they had considered programming in Forth, they expressed the opinion that Forth is no match for C.

Are they right?

Besides, B.E.S.T. says they have developed their own auxiliary C tools that allow them to program on the Amiga twenty times faster than ordinary Aztec C.

There are lots of libraries of subroutines which can be purchased and incorporated into C programs to speed development time. Are there libraries of extra Forth "words" for sale to speed development? It seems to me that if Forth has been pushed aside by C, it may be because brilliant Forth programmers are spending their time reinventing the wheel (laboriously adding words to Forth and fiddling with how things are going on and coming off the stack) while C programmers are building upon the labors of those who have already invented the wheel, and they are making better progress.

Or, put another way, C is so difficult to learn and use that those who do use C are willing to pay for and include other programmers' work in their own programs. Forth, on the other hand, is so much easier to learn and use that Forth programmers have fallen behind while they happily work at adding their own words to extend their personal copy of the language.

I read in the August 1985 Computer Language that Philippe Kahn says Bordeaux will offer Modula-2, BASIC, C and Ada languages, but not Forth. He won't offer Forth because, he says, "Forth is a religion . . ."

Where can I read the definitive comparison of C and Forth? Are Forth programmers sharing words instead of encouraging Forth users to go off on their own to reinvent the wheel?

Awaiting your prompt reply,

Rich Kevin O'Brien 
Renton, Washington

"Forth as a religion" was a way of poking fun at our own conviction, but like many analogies it has both detracted from real issues and outlived its usefulness. One real issue in need of more attention is Forth "subroutine" libraries. John James, in particular, has studied the factors involved in providing in individual modules the useful tools developers want and need. One assumes that a stable Forth nucleus will be critical to the creation of such off-the-shelf, plug-in extensions. Both system vendors and third parties may find a line of profit in this thinking, provided the work meets specific design criteria (see James' articles in Forth Dimensions VII/4, VII/5 and VIII/2).

By the way, Dr. Dobb's Journal still publishes an annual Forth issue and, as of this writing, a periodical column about Forth. I haven't seen any "definitive comparison of C and Forth," but if one or more authoritative, objective authors could be found I'm sure the editors of DDJ and those of Computer Language would be interested in the idea. But I was talking with CL Publications' co-founder Craig LaGrow about this when he made the interesting point that the worst published comparisons are those that attempt to compare apples and oranges. Perhaps that is the area that will have to be addressed by anyone comparing Forth and C.

—Editor

Teaching Forth:
Testing TESTIT

Dear Marlin:

Mr. Apra's article, "Let's Keep It Simple" (VII/6) makes a good point, though perhaps not strongly enough. I have done some teaching and can testify to the real value of simplicity. Nevertheless, I would like to take exception to the efficacy of TESTIT, IFTRUE, IFFALSE and ENDIT as Mr. Apra proposes their use.

The points I would make are (refer to the Apra screens):

1. Students ought to learn how to recognize the need for DUP and DROP. It may be confusing that TESTIT will sometimes precede and at other times follow the test depending on whether the number is to be tested or the resulting flag needs to be DUPed. In the definition of EXAMPLE1 a DROP would need to be added if one of the two phrases were eliminated. It would be better to replace TESTIT with DUP. Then the correct action of eliminating DUP rather than adding DROP would be evident.

2. It is important that one's understanding of the IFFALSE ELSE ENDIT control structure be easily transferable to the IF ELSE THEN constructs. The major problem here, as I see it, is demonstrated by the definition of EXAMPLE4, which has the appearance of a CASE statement. I can imagine it might lead to further confusion. Of course, with a CASE statement, the ending DROP would be unnecessary, as well as the numerous TESTITs (DUPs).

3. Mr. Apra correctly observes that "... the IF ELSE THEN syntax does not leave enough clues to where the parts of the conditional should go." But in the proposed syntax there is no clue to suggest when — and when not — to include TESTIT. Compare the definitions of EXAMPLE1 and EXAMPLE2.
I offer the code on the enclosed two screens as the second step in the quest to simplify IF ELSE THEN, for the purpose of teaching, to an understanding of that conditional structure. I know from experience that such conditionals are difficult to grasp for many a beginner. Sometimes those of us who do understand forget that.

I believe the IFTRUE OTHERWISE RESUME concept offers a better chance to gain an understanding of the conditional’s true structure. The examples I have included were designed to demonstrate how variable the structure of such conditionals can be. Replacing DUP with TESTIT in the examples will show why I think TESTIT is not so useful.

I offer my congratulations to Mr. Apra. I hope there are many like him searching for better ways to teach beginning young programmers. Finally, a word to the students: letters such as this one are not “put downs” of other authors. Progress doesn’t stop with a first attempt at improvement; it is an evolution. I am confident that competent programmers will argue with both Mr. Apra and me. Some will even poo-poo it altogether. So be it. But, of course, most of them are not teachers.

Sincerely,

Gene Thomas
Little Rock, Arkansas

Listing 1
Screen #20
0. IFTRUE/IFFALSE, OTHERWISE (optional), RESUME GT May86
1. IFTRUE \ f -- (if if execute following)
2. [COMPILE] IF \ IMMEDIATE
3. IFFALSE \ f - (if ff execute following)
4. COMPILE NOT [COMPILE] IF \ IMMEDIATE
5. OTHERWISE \ execute following if IFTRUE or IFFALSE not exec.
6. [COMPILE] ELSE \ IMMEDIATE
7. RESUME \ -- (resume program here after IFTRUE-OTHERWISE)
8. [COMPILE] THEN \ IMMEDIATE
9.
10. ?EQUALSTEN \ n -- ) DUP 10 =
11. IFFalse." True =10\ DROP OTHERWISE ." False " 10 >
12. CASE 0 OF ." <10" ENDOF
13. 1 OF ." >10" ENDOF
14. ENDCASE
15. RESUME ;
16. IFTRUE; scan #2; examples
17.
18. ?EQUALSTEN \ n -- ) DUP 10 = DUN
19. IFFalse." Yes," SWAP , RESUME
20. IFFALSE ." No," . RESUME ;
21.
22.
23. ?EQUALSTEN \ n -- ) DUP 10 =
24. IFFalse." True" DROP ( and proceeds to last resume)
25. OTHERWISE 10 > (compare with \n )
26. IFFalse ." Bigger than 10"
27. OTHERWISE ." Smaller than 10"
28. RESUME
29. RESUME ;
30.
31.

FORTH Dimensions 6 Volume VIII, No. 3
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Return to Forth Interest Group, P. O. Box 8231, San Jose, CA 95155 • 408 277-0668
When New Becomes Obsolete

Dear Sir:

Stop, stop, for pity’s sake, STOP! In anguished supplication I beg of you, whoever you may be: Please, not an 87-Standard Forth! The way things are going, in a few years you will have to change the name *Forth Dimensions* to “*Forth Fragmentations!*”

I concede that Forth-83 is not perfect, nor was fig-FORTH or Forth-79. But the situation is getting ridiculous with so many “official” versions of Forth. Is this the planned model change every four years that once made us love GM? I submit that no one will benefit from any further changes in “standard” Forth except those who would sell new Forth systems to the few unsuspecting but trusting souls who believe they must replace an “obsolete” system.

fig-FORTH was never so bad that it had to be junked; it is my firm conviction that [the change to] Forth-79 was cosmetic, and Forth-83 even more so. Why make it so difficult for fig-FORTH users to talk to Forth-79 or Forth-83 users? I think a much better solution to the problem of deficiencies in fig-FORTH would have been to develop a standard set of extension words for fig-FORTH which would have been added to existing systems at the discretion and convenience of the programmer.

It is still not too late! If we must live with the present standard Forth-83, then so be it. But let us not compound the felony by foisting another “standard” on the already troubled Forth community.

Yes, I do have an operating Forth-83 system. It is the excellent offering by Wilson Federici for the 6809; he is giving it away! However, I am seriously considering going back to my 6809 fig-FORTH, which I have been working on to use bank switching something like the 8088. I realize I can’t make fig-FORTH behave exactly like Forth-83 by just adding additional words, but I hope to get close. However, I do know this — I will never update to a new standard Forth!

Let me address another point as related to standards. I have seen several letters lately asking for changes in Forth that would accommodate the thirty-two-bit crowd with their larger addressing space. Don’t change Forth! Instead, add some new words which will process the extra addressing. What is wrong with words like 4DUP, etc.? Additions like this would still make it possible for people like me, who plan to keep their eight- and sixteen-bit machines (like my 6809), to continue talking to the rest of the Forth community.

Perhaps I am in the wrong, though. Perhaps Forth is primarily used by professional programmers and system designers who simply tolerate the hobbyist crowd as long as they keep quiet. Am I a voice in the wilderness, talking only to myself? Am I the only hobbyist who programs in Forth because it is his favorite language? Oh, I have written

---

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commercial programs for engineering and control, but in other languages (BASIC and Fortran) because the customer was afraid of a language “nobody ever heard of.” Maybe a few more hobbyist users of Forth would actually make it easier to sell Forth programs.

Am I talking to the wrong group? Possibly FIG, too, is not interested in the hobbyist. Am I wasting my money belonging to an organization which appears to treat me with disdain? The questions are not strictly rhetorical; I hope someone will answer me. My concern about proliferating standards is related to program exchange among hobbyists. Already there is a problem with IBM? More “standards” will only make matters worse.

I would like to close with one more question, “Who is kidding who?”

Sincerely,
R.D. (Doug) Lurie
Leominster, Massachusetts

Terminal TI

Dear Editor,

As a relatively recent convert to Forth (little more than a year of Forth), I was delighted to find that your publication is actively searching for articles on TI-Forth for the 99/4A. I write a column in our user-group newsletter (MSP99) on Forth.

With reference to “Simple Data Transfer Protocol” (Forth Dimensions VI/2), I submit the enclosed two pages. The first, second and last screens are my adaptation of the protocol to the 99/4A. The terminal emulator will work in half or full duplex. The routines are set up to use RS-232 port one at 300 baud. Data for alterations is given in the comments. Speeds greater than 1200 bps may require some routines to be in CODE. The only option required is the “-CRU” option, loaded in the first screen (that’s what the “CLOAD” loads).

Cordially,
Glenn Davis
St. Paul, Minnesota

---

```forth
(serial communication constants TI 99/4A 26SEP85 GED)
BASE->R DECIMAL 88 R->BASE CLOAD STCR
BASE->R HEX

0408 CONSTANT RATE (0400=300 bps; 0410=1200 bps)
003 CONSTANT PROTOCOL ( 0 bits, no parity, 1 stop)
01300 2 / CONSTANT CARD (1300=ports 1&2 1500=ports 3&4)
CARD 07 + CONSTANT LED (LED on RS232 card)
CARD 02B + CONSTANT PORT DECIMAL (02B=odd port 040=even)
PORT 13 + CONSTANT LDIR (Load Interval Register)
PORT 16 + CONSTANT RTSON (Request To Send ON)
PORT 21 + CONSTANT RIENB (Receive Interrupt ENABLE)
PORT 22 + CONSTANT XBRE (transmit Buffer Reg Empty)
PORT 31 + CONSTANT RESET (RESET TMS9902 ACC chip)

(serial communication primitives TI 99/4A 13JUN85 GED)
BASE->R HEX

XINIT 00 00 10 1F 00 00 00 00 (Reset serial port)
PROTOCOL 08 PORT LDCR (load rate, parity, stop bits)
LDIR SBZ (inhibit access to interval reg)
RATE OC PORT LDCR (load rcv, xmt rate registers)
?XOUT (---- f) XBRE TB
?XIN (---- f) RBRL TB
?XOUT (char ----)
XOUT 00 PORT STCR RIONB SBZ LED SBZ
XIN (--- char)
LED SBZ BEGIN ?XIN UNTIL 08 PORT STCR RIONB SBZ LED SBZ

(communication protocol primitives TI 99/4A 13JUN85 GED)
BASE->R HEX

TXCLR 0 XOUT;
END
BEGIN 5 (enq) XOUT XIN 6 (ack) = UNTIL;
ACK
BEGIN XIN 5 (enq) = UNTIL
6 (ack) XOUT;
2DUP
OVER OVER
DROP;
DROP DROP;

(transfer protocols primitives TI 99/4A 13JUN85 GED)
BASE->R DECIMAL
SEND-BLOCK (block-addr ----)
DUP B/BUF + SWAP
BEGIN DUP 00 XOUT 1+ 2DUP =
UNTIL 2DROP;
TAKE-BLOCK (block-addr ----)
DUP B/BUF + SWAP
BEGIN XIN OVER C! 1+ 2DUP =
UNTIL 2DROP UPDATE;
SYNC TXCLR
BEGIN 5 XOUT
?XIN IF XIN 5 = ELSE 0 THEN
UNTIL 6 XOUT
BEGIN XIN 6 = UNTIL;

R->BASE -->
```

---

(continued from page 9)

```forth
(continued from page 9)

```

---

(continued from page 9)

```forth
(continued from page 9)

```
Multiple LEAVEs by Relay

Dear Editor:

John Hayes’ “Another Forth-83 LEAVE” (VI/1) stimulated me to try to find an even simpler way to handle multiple Forth-83 LEAVEs. I decided that a straightforward approach involved having each LEAVE simply branch to the next LEAVE, with the last one removing the index values from the return stack and branching to the word following LOOP.

This is accomplished by using a flag to show if there is a LEAVE in the loop; if so, LOOP sets up a forward branch. Also, if there is already a LEAVE present, an added LEAVE sets up a forward branch to the new LEAVE, not to the next word. Thus, when LOOP is reached, there cannot be more than one LEAVE whose branch jump needs resolution. In the code, (DO) and (LOOP) are the standard constructions, as in McCabe’s Forth Fundamentals.

Sincerely,

Chester H. Page
Silver Spring, Maryland

```
(transfer protocol user lexicons TI 99/4A 13JU85 BAD
based on words from Forth Dimensions VI/2) BASE->R DECIMAL
XMT (first-scr# last-scr# ---)
SYNC 1+ SWAP
DO ENQ 2 XOUT 1 BLOCK SEND-BLOCK
LOOP ENQ 4 (bot) XOUT;
RCV (first-scr# ---)
SYNC
BEGIN DUP BLOCK ACK XIN 2 (stx) OVER =
IF DROP TAKE-BLOCK 0
ELSE 4 (bot) =
IF DROP FLUSH
THEN 1
THEN SWAP 1+ SWAP
UNTIL DROP;
R->BASE -->

(TI 99/4A terminal emulator) 26SEP85 BAD
BASE->R HEX
0 VARIABLE ECHO
: TERMINAL XINIT
BEGIN
?XIN IF XIN EMIT ENDIF
?KEY -DUP IF
DUP 2 = IF QUIT ENDIF
0B37C @ B2B AND
IF
ECHO @ IF DUP EMIT ENDIF
XOUT ELSE DROP
ENDIF
AGAIN;
R->BASE;

)```

```
(LEAVE) R> R> DROP DROP BRANCH;
: >RESOLVE ( addr---) HERE OVER - SWAP ! ;
: <RESOLVE ( addr---) HERE - , ;
0 VARIABLE LEAVE.FLAG
: DO 0 LEAVE.FLAG ! COMPIL (DO) HERE 3 ; IMMEDIATE
: LEAVE LEAVE.FLAG @ IF >RESOLVE THEN LEAVE.FLAG !
COMPIL (LEAVE) HERE 0 ; IMMEDIATE
: LOOP 3 ?PAIRS COMPIL (LOOP) LEAVE.FLAG @
IF >RESOLVE THEN <RESOLVE ; IMMEDIATE
```
Hackles and Hopes

Anyone who has read announcements of this year’s FORML conference undoubtedly recalls the theme, “Extending Forth Towards the 87-Standard.” Already this has raised both hackles and hopes among those who care about standardization work.

Some companies with huge Forth-79 programs chose not to switch to Forth-83 because of the in-house costs of upgrading functional systems. Some vendors felt it more important to provide a stable system to their existing clientele than to incorporate the changes required by Forth-83. And many assert that redefining words in the nucleus is unforgivable and potentially lethal to Forth’s commercial viability.

How we view the Forth standard depends on who we are. System vendors have a commercial interest in product stability that hobbyists don’t share. Professional users welcome functional improvements if they are of clear benefit, but resist energy-consuming changes unrelated to getting a good job done on time. Some language hackers embrace any opportunity to save a clock cycle and will modify their personal systems regularly without qualm.

The interests of its diverse users brought them to Forth in the first place, and meeting their needs perpetuates the language’s use. Those needs are as real and specific as the measurable performance of a primitive word and must, as a matter of course, be part of the overall approach to the standards effort. It is evident that the Forth community has many interests, of which the most fundamental is success of the language itself. A standard arises from that common interest and does resist to changes. The intensity and volume of debate over Forth-83 obscured that common interest, creating an us/them, win/lose atmosphere. I hope anyone participating in future debates will have studied Getting to Yes by Roger Fisher and William Ury of the Harvard Negotiation Project.

This morning I spoke with Guy Kelly, chairman of the Forth Standards Team (FST), about the growing discussion over another Forth standard. He told me the FST has no plans for a Forth-87 and has not called any meetings on the topic. During two meetings leading to the Forth-83 Standard, the team had to deal with about five hundred proposals; at that time, members of the team decided there would be no future FST events without advance technical meetings at the regional level. Proposals would then be published and distributed to subscribers. Guy said that, as of his last poll, FST referees had not indicated any desire to begin this process. It is easy to see that, even if a new standard is proposed, 1987 is an unlikely target date for its completion.

Forth Dimensions maintains a neutral ground for discussion and accepts articles and letters related to this subject. We hope representatives of the standards team will write about its current plans and intentions, as well as the focus and likely restrictions of future FST meetings. Users and vendors are equally welcome to submit their viewpoints and suggestions. Our pages are limited, so items related to general topics (e.g., extension word sets, upward compatibility), overall direction and philosophy will carry the greatest weight; most experimental, technical proposals would be better placed at FORML. Forth Dimensions will attempt to take no sides, warning only that we are looking for more light than heat. Coherent discussion is welcome from all “sides” in the debate.

The proper place to send formal proposals is the Forth Standards Team (P.O. Box 4545, Mountain View, California 94040). A proposal (even one suggesting only that the Forth kernel remain unchanged) must be submitted in the manner established in the published Forth-83 Standard document, and which we will reproduce in these pages if space permits. Copies of the standard are available from several vendors and the Forth Interest Group.

As always, we will continue publishing material that should be instructive and useful to our readers right now. Enjoy the issue!

—Marlin Ouverson
Editor
Forth Systems With a

Segmented Memory Model

Richard Wilton
Marina del Rey, California

Most Forth interpreters confine themselves to a single logical address space in memory. Executable code, linked lists of addresses, and structures of initialized and uninitialized data are all stored within the same address space. Accessing any of these program elements within a Forth program is simple because any executable code or data item can be reached with a unique sixteen-bit address.

The lack of a structured "memory map" in Forth systems has not proven to be a major deficit in small computers with sixty-four Kbytes or less of main memory, or in computers without sophisticated operating systems. However, in complex microcomputer operating system environments, the use of a segmented memory model within a Forth interpreter offers clear advantages over older approaches.

Advantages of Memory Segmentation

The Intel 8086. Initially, the prevalence of computers based upon the Intel 8086 family of microprocessors provided the impetus for creating a Forth interpreter with a segmented memory model. The architecture of these processors is such that memory is most conveniently addressed in sixty-four Kbyte partitions or segments. Since the 8086 and related processors have four segment addressing registers, a software memory model which "maps" into these registers makes very effective use of memory.

For example, a memory model which uses four non-overlapping segments can conveniently utilize 256 Kbytes of RAM in an 8086-based machine, even though the software handles only sixteen-bit addresses for the most part. A Forth interpreter can thus make effective use of up to 256 Kbytes of RAM even with a sixteen-bit address interpreter.

Headerless and ROMable Code. There are other strong reasons for segmentation of memory in a Forth system apart from such hardware-related considerations. For example, applications developers who have no need for dictionary headers can easily excise them from a finished application if the headers are maintained in a separate memory partition. Also, the separation of executable code from other data greatly simplifies the implementation of a Forth interpreter which can be executed from read-only memory (ROM).

Links to High-Level Languages. Still another important reason for using a Forth interpreter with a segmented memory model is that most high-level language compilers generate executable modules with similar memory segmentation. It is easier to link to such "external" modules from a Forth system with a segmented memory structure.

The Basic Segmented Memory Model

The implementation of a Forth system with segmented memory is straightforward. The basic memory map consists of three discrete partitions or segments:

- Code segment — contains all executable code
- Data segment — contains bodies of colon definitions as well as initialized and uninitialized program data
- Headers segment — contains names and addresses of all definitions

The parameter and return stacks may be placed in a fourth discrete segment. However, it is simpler in many respects to maintain the stacks in one of the other segments. Furthermore, certain operating systems constrain the stack to either the code or data segments.

Theoretically, there is no restriction on the amount of segmentation that can be done in creating the memory map. For instance, disk buffers, constants, or a "user" area for multitasking might each be mapped into separate segments of their own. However, the increased overhead involved in keeping track of many separate segments can degrade system performance as well as increase the complexity of the system.

The separation of executable code from data is particularly useful in environments in which memory protection schemes are used. For example, in its "protected" mode the Intel 80286 supports the creation of execute-only code segments and non-executable data segments; attempts to execute data or to modify existing code generate hardware exceptions which can trap to software error-handling routines. This feature might be much appreciated by Forth programmers whose programs would otherwise crash irretrievably due to such errors.

System Data Structures

The structures which contain the data upon which the interpretive Forth system operates reflect the segmentation of the system's memory map. The actual data structures used in a particular implementation necessarily depend upon the type of threaded code used.

Indirect-Threaded Code. For example, in an Indirect-Threaded Code (ITC) system, a colon definition is mapped into a list of sixteen-bit addresses in the data segment, as in Figure One.*

The code field contains the address of the inner interpreter's "nest" routine. Since this field points to executable code, the address refers to a location in the code segment. The body of the definition contains code field addresses (compilation addresses), all of which refer to locations in the data segment.

A CODE definition is mapped primarily into the code segment. The only part of a CODE definition to be found in the data segment is the address of the executable code, which is stored in the data segment at the definition's compilation address. (See Figure Two.)

Direct-Threaded Code. The structures are somewhat different in a Direct-Threaded Code (DTC) system. In a DTC implementation each colon definition references the inner interpreter's "nest" routine explicitly with a short fragment of executable code, as shown in Figure Three.
This executable fragment (a jump instruction) resides in the code segment. The "nest" routine uses the address stored in the code segment to find the body of the definition, which resides in the data segment. Similarly, a CODE definition resides entirely in the code segment.

Headers. The actual structure of dictionary headers, which are maintained in a segment of their own, greatly depends upon the structures used to represent Forth vocabularys. However, the basic structure is simple, as in Figure Four.

The "compilation address" field deserves careful attention because the type of address it contains depends upon the implementation. In an ITC implementation, this address references a location in the data segment. In a DTC system, the address is located in the code segment. This difference is critical in understanding the function of a Forth system with a segmented memory map, although it is immaterial to a Forth applications programmer. (For example, the sequence ' <name> execute has the same effect in either implementation.)

Implications for System Implementation

Because code and data are separated, complex and unconventional colon definitions can be easily built. For example, the implementation of words containing create ... does or ;code is very straightforward, especially in a DTC system where the size of the code fragment associated with a definition is not limited.

Since compilation addresses are maintained explicitly in dictionary headers, it is easy to implement a forwarding-referencing scheme. A forward reference is created with a "dummy" compilation address. Later, when the reference is resolved, the "dummy" address is replaced by a real one.

Both Unix and MS-DOS offer dynamic memory allocation facilities. Such dynamically-allocated memory is usually located in a partition outside of a requesting program's initial address space. A Forth which integrates a segmented memory model is easily adaptable to this scheme of memory allocation.

In Forth systems with partitioned memory, Forth memory operators such as @ and # implicitly address the data segment. A problem arises when it is necessary to access other memory partitions explicitly. One simple approach is to utilize "long" memory operators which require both a "base address" (e.g., a segment address) and an "offset":

@L ( base_addr offset -- value )
#L ( value base_addr offset -- )

CMOVEL ( base_addr1 offset1 base_addr2 offset2 #bytes -- )

This approach works well in practice. In any case, such "long" operations comprise only a small percentage of most applications.

The unique advantages of a Forth interpreter with a segmented memory model are not restricted to microprocessors which directly support memory segmentation. The segmented memory scheme works very well on the Motorola 68000, for example. Of course, for older eight-bit processors with only 64 Kbytes of address space, this sort of memory partitioning involves more bookkeeping overhead than it is worth. Nevertheless, a partitioned memory model in Z-80 or 6502 systems with bank-switched memory, such as the Apple IIe, may yet prove to be useful.

The performance of a Forth interpreter with a segmented memory map is not significantly affected by the separation of code, data, and headers. Other factors, including the method of dictionary search, the nature of the threaded code implementation (e.g., ITC or DTC) and the efficiency of machine code primitives, influence performance to a much greater degree.

For example, execution speed benchmarks on an ITC implementation of an 8086-based Forth system (PC/Forth by Laboratory Microsystems) are virtually identical whether or not the memory map is segmented. Considering the improvement in efficiency and flexibility of memory utilization, a Forth system which incorporates a partitioned memory model is a better approach to Forth programming in demanding sixteen-bit systems with complex software environments.

*See figures on page 27.

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The Point Editor

J. Brooks Breeden
Columbus, Ohio

Hardin's Law: You can never do just one thing.

Often, I find books and back issues of *Forth Dimensions* contain tips and techniques that I didn’t understand or need when I first read them but that now make perfect sense and are exactly what I’m looking for. Sometimes, however, what I need to help me solve a problem is hiding, disguised as a solution to a different kind of problem. For example:

Recently, while developing a computer-assisted instruction (CAI) drill-and-practice test on storm drainage, I wanted to draw a map of the United States, showing zones of approximately equal rainfall. Because I wanted just one small map in a corner of the display, I found an 8x10-inch map of the USA and traced it onto graph paper, redrawing the outline with straight line segments that gave a reasonable facsimile of the USA's shape. Next, by trial and error, I found a scale which related the USA's vertical dimension to the y-axis coordinate range I wanted on the display, and scaled the x,y coordinates based on 0,0 in the upper-left corner of Washington state. The USA outline required the first 128 x,y coordinate pairs shown in Figure One. The remainder of the points in Figure One are used to draw the lines of “approximately equal rainfall.”

I use PC/Forth 3.1 from Laboratory Microsystems, Inc. and, normally, CAI graphics are simple to program using the supplied routines for LINE, FLOOD, ARC, etc. I was worried about running out of memory with this application, however, and I didn’t want to compile a word consisting of $x \ y $ LINE 128 times for the USA’s outline! I had also measured the x and y coordinates using the same scale, and the display required an x-axis multiplier because the pixels on IBM displays aren't square. Having gotten this far, I finally stopped to think about the problem. How was I going to draw the USA?

Three methods came to mind:

1. Define a command to DRAW from the last point to the next point, letting the computer re-scale the x-axis and add an offset from an “x,y origin” on the display. I could use a brute-force approach with screens full of \texttt{X Y DRAW} instead of \texttt{X Y X Y LINE} (after all, it was just one map), and I could relocated the map on the display by changing the coordinates of the origin.

2. Create the \texttt{DRAW} command as above, manually re-scale the x-axis and (comma) the coordinates into an array. A loop could read the array and draw the USA’s outline.

3. Recall a previously stored image from disk with PC/Forth’s \texttt{GET-IMAGE} and overlay it with the rest of the tutorial.

The first method would require several screens of \texttt{X Y DRAW} and matching a point would require looking through literally hundreds of numbers and \texttt{DRAWS}. The second method would take fewer screens, but editing numbers separated by only commas (no \texttt{DRAWS}) would be even harder. Using (comma) to lay down all the points in the dictionary would use lots of memory, too. The third method wouldn’t solve anything, since I had to draw the image, somehow, before I could store it!

BASIC draws points from arrays stored in disk files all the time! Then, I remembered an exercise in *Starting Forth* (chapter ten, exercise five, page 288) to write words to accept pairs of numbers from the keyboard, write them to a block, then recall and format them on the display. Eureka! I could store coordinate pairs in a block and draw the USA’s outline by using a loop to read coordinates from the block.

Hardin’s Law Strikes

Have you ever entered 256 consecutive numbers? Did you ever make a mistake? Don’t answer! It’s not enough to just write numbers to a block and read them back. You have to be able to change a number and you need to be able to add or delete one, or several (in the middle, yet). You see, you can never do just one thing! I needed a “point editor.”
POINTED.SCR

POINTED ("Point Ed-itor") stores points in a block. Using a non-DOS, block-oriented FORTH like polyFORTH, that's it; you have a block of points. If, however, your Forth uses screens in a DOS file, will the point data be a screen in your application's file, or a separate file on the disk? Because I wanted to be able to read the point data from a PC/Forth TURKEY application (a DOS .COM file which wouldn't have any blocks), I needed to store the points in a separate, named, "point file" on the disk. The DOS interface was something I hadn't intended to deal with: how and when to open and close screen files within Forth, from a .COM file, FCBS, handles, etc. Hardin's Law.

To put off dealing with DOS for awhile, I chose to use POINTED's screen ten, as a data block for "editing." I entered all the required points and used LMI's SCOPY utility both to create a new file and to transfer POINTED's screen ten to the new file's screen zero. (Screen zero can be read and written to, but cannot be loaded; a distinct advantage)

<table>
<thead>
<tr>
<th>Points list for screen 10</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pt.</td>
</tr>
<tr>
<td>----</td>
</tr>
<tr>
<td>0</td>
</tr>
<tr>
<td>5</td>
</tr>
<tr>
<td>10</td>
</tr>
<tr>
<td>15</td>
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<tr>
<td>160</td>
</tr>
<tr>
<td>165</td>
</tr>
<tr>
<td>170</td>
</tr>
</tbody>
</table>

Screws on one through four contain the "command-driven" program. Defined words allow scaling the x-axis, entering, changing, deleting, listing and printing points stored in block ten. Screens five through seven make a simple menu-driven interface that a graduate student assistant not familiar with Forth can use to input points (Hardin's Law ... Murphy's Law, too). I like to think the source for these screens is self documenting.

Screen eight contains words for re-scaling between "raw" data and CGA/EGA scaled data. These words aren't normally loaded, as they can be dangerous, although I suppose the most dangerous word is an option on the menu. The "best" scale factor for CGA/EGA will vary with the vertical adjustment of your display. (Yours will probably be less.) Experiment, but backup block ten before you mess around — it can be frustrating to lose the original data and have to re-enter it.

A more sophisticated version would create a file and work directly with its screen zero without POINTED's screen ten as the "middle man." Another improvement would automatically up-
date the number of points in the file, eliminating the hassle of manually keeping track of the number of points. Still, if you store sets of points to draw different things in the same block, such as the “lines of approximately equal rainfall,” you will have to keep track of starting and ending point numbers anyway. There are many other improvements to POINTED that I leave to the reader to implement; but remember, I wasn’t trying to write a universal point editor. I was trying to draw just one map!

Screen nine contains the drawing routines. This screen is not a part of POINTED.SCR, per se, but is included in the file for convenience. Normally, you copy it to your application and modify the file names, origin, etc., as required. These definitions should be self-explanatory with the exception of V+ which adds vectors. In DRAW-PTS, PT# fetches a point’s coordinate pair and the ORIGIN’s coordinate pair. V+ adds the ORIGIN’s x to the point’s x and the ORIGIN’s y to the point’s y. By changing the origin, you can draw the same image at different places on the screen. This allows “fine tuning” the display. Note that the file names compiled by DRAWER and POINTFILE have extensions. If you TURNKEY the application, be sure the extension for the name in DRAWER is .COM.

Listing One

Screen # 0

( POINTED Point Editor for drawings jbb 10:45 04/28/86 )
( Last change: Screen 002 jbb 15:31 04/29/86 )

Copyright 1986, by J. Brooks Breeden
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You may freely use this program for personal non-commercial use.

The point editor creates files of binary points (coordinates) stored in Screen 10 of THIS file. Up to 224 PAIRS of x y values may be stored. Use SCOPY to create a new file and transfer data in Screen 10 to the new file’s Screen 0.

DRAWER.SCR shows how to call the points from another program.

This program is written in and for the Forth-83 implementation PC/FORTH 3.1 by Laboratory Microsystems, Inc.

Screen # 2

( Enter points jbb 15:31 04/29/86 )
: ENTER-PT ( - x y ) CR â€ “ Enter x <CR> y <CR>. ” CR
. . ” X = ” BIN CR . . ” Y = ” BIN CR ;
: WAIT 224 GOTOXY . ” Press any key to continue.” KEY DROP ;
: ?PTS ( - start count ) . ” Starting at which point # ? ”
CR BIN 224 MIN CR . ” How many points? ” BIN 224 MIN ;
: ?-PTS ( - start end ) . ” Delete from point # ? ” CR BIN
224 MIN CR . ” Delete thru point # ? ” CR BIN 224 MIN ;

Screen # 3

( Change & delete pts. jbb 15:41 04/27/86 )
: PTS ( start count - ) SWAP POINTSBLOCK SWAP 2* 2* DUP >R
OVER + R> CMOVE UPDATE ;
: -PTS ( 1st last - ) 1+ 224 OVER - 2* 2* +R
POINTSBLOCK SWAP POINTSBLOCK R> CMOVE UPDATE ;
: CHANGE-PT CLS SCALEIT? CR . ” Enter the point # ” CR
IN ENTER-PT SCALEING-x? ROT IPOINTS UPDATE SAVE-BUFFERS ;
: ADDPTS ( start count - ) CLS SCALEIT? CR OVER + SWAP
DO . ” Pt. # ” 1 . , ENTER-PT SCALEING-x? !POINTS LOOP ;
: INSERTPTS CLS ?-PTS 2WP +PTS ADDPTS SAVE-BUFFERS ;
: DELETPTS CLS ?-PTS -PTS SAVE-BUFFERS ;

Screen # 4

( Show & print pts jbb 10:46 04/28/86 )
: .PTS ( start count - ) . ” Points list for screen 10” CR
. ” Pt. # X Y ” OVER + SWAP
DO 0 I POINTSBLOCK 2* SWAP I 5 MOD 0=
IF CR ELSE 3 SPACES THEN
. ” # I 3 U.R SPACE 4 U.R 4 U.R LOOP CR ;
: SHOWPTS CLS ?-PTS CLS WAIT ;
: PRINTPTS CLS ?-PTS PRINTER .PTS CONSOLE ;
: WIPE-BLOCK-10 ( - ) POINTS BLOCK 1024 BLANK ;
( DANGER!!! If you do this you wipe it ALL away!! )

Screen # 5

( Choices jbb 10:47 04/28/86 )
: CHOICES BLUE BG CLS WHITE FG 35 7 GOTOXY . ” MAIN MENU”
GRAY FG
32 11 GOTOXY . ” (S)how points”
32 12 GOTOXY . ” (I)nter points”
32 13 GOTOXY . “ C)hange a point”
32 14 GOTOXY . ” (D)elete points”
32 15 GOTOXY . ” (P)rint points”
32 16 GOTOXY . ” (M)ove block 10!!! (dangerous!!!)”
32 18 GOTOXY . ” Esc leaves the program” WHITE FG
26 20 GOTOXY . ” Press the first letter of your choice.”
GRAY FG ;
The second listing, DRAWER.SCR, is a demo to show that DRAWER does indeed draw points entered with POINTED and that the result does resemble the USA. Screens one through eight build a dummy of the application that started the whole thing. An EGA screen dump of the display is shown in Figure Two. It's ironic that the source to dummy the application is almost longer than everything else. I could not avoid the DOS interface forever. Screen five shows how I sequenced the opening and closing of files so the program doesn't get lost in DOS and crash, leaving files lying around open.

The application uses EGA graphics and floating-point math extensions. The demo dummy lacks about thirty screens of irrelevant, generic CAI question drivers, answer checkers and random problem generators. Everything looks exactly like the real McCoy, except it doesn't work. You do not need floating-point math routines to load and run the demo, but you do need an EGA card and display to run the demo as it is written. (Attempting to run the demo from PC/Forth without loading EGAGRAPH first will result in a truly spectacular hard crash!) To run with a CGA display, eliminate color references, re-scale the points and adjust all y-axis references in DRAWER.SCR accordingly.

I hope some of the "other things" I had to do in trying to draw just one map may be of help to others trying to do just one thing, sometime. If anyone is interested in the inner workings of the real application, please contact me. I'd be happy to discuss the specifics.

---

Screen # 6
(Message jbb 10:47 04/28/86 )
: MESSAGE CLS 10 GOTOXY
" Note: Points will be stored in Screen 10 of this file."
29 14 GOTOXY " Be certain Screen 10 is wiped or that"
29 15 GOTOXY " it is the point-set you wish to edit."
29 16 GOTOXY " Use SCOPY to make a new file and to copy"
29 17 GOTOXY " Screen 10 to Screen 0 of the new file."
-->

Screen # 7
(Main program jbb 10:47 04/28/86 )
: MENU CHOICES KEY
DUP 2 T = IF DROP EXIT ELSE
DUP ASCII S = IF DROP SHOWPTS ELSE
DUP ASCII I = IF DROP INSERTPTS ELSE
DUP ASCII C = IF DROP CHANGE-PT ELSE
DUP ASCII N = IF DROP DELETEPTS ELSE
DUP ASCII P = IF DROP PRINTPTS ELSE
DUP ASCII W = IF DROP WIPE-BLOCK-10 ELSE
DROP 50 50 BEEP THEN
THEN THEN THEN THEN THEN RECURSE ;

Screen # 8
(Optional dangerous words jbb 15:24 04/28/86 )
\ There is obviously rounding error in these words.
: UNSCALE-EGA (-) 173 0 ( convert EGA scaling to raw data)
DO I POINTSBLOCK 20 SWAP 12 */ SWAP I !POINTS
LOOP ;

Screen # 9
(Drawer routines jbb 15:30 04/29/86 )
\ Copy this block into your application and modify as needed.
VARIABLE CP ( holds coordinates of graphic cursor)
DO 1 POINTSBLOCK 20 CP 13 */ SWAP I !POINTS
LOOP ;

Continued on page 23.
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Listing Two

Screen # 0
(Drawer EGA LMI PC/Forth version jbb 15:13 04/20/86)
(Last change: Screen 002 jbb 16:25 04/28/86)

This file requires an EGA adapter and display. You must load
PC/FORTH's EGAGRAPH.COM before loading PC/FORTH in order to
load this file, or you will CRASH!!! I guarantee it!

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This program is written in and for the Forth-83 implementation
PC/FORTH 3.1 by Laboratory Microsystems, Inc.

---

Screen # 1
(Utilities required for demo jbb 15:20 04/20/86)

: FG FOREGROUND ; (shorthand notation for PC/FORTH's names)
: BS BACKGROUND ;
: HUE (color-) CREATE, DOES> (-n) ;
7 HUE GRAY 15 HUE WHITE 1 HUE BLUE
4 HUE RED 2 HUE GREEN

2VARIABLE CP
(holds coordinates of graphic cursor)
: AT (x y -) CP 2! ;
: DRAW (x y -) 2DUP CP 20 LINE AT ; (draw line, update CP)

---

Screen # 2
(Drawer routines jbb 16:25 04/29/86)

: DRAWER "DRAWER.SCR" ;
: POINTFILE "USAEGA.SCR" ;

0 CONSTANT POINTS
380 66 2CONSTANT ORIGIN

\ Note: the POINTS defined in POINTED.SCR is not the same.

POINTS (n-) POINTS BLOCK + 64 + ;

: PTX (n - x y) 2* 2* POINTS BLOCK 20 ORIGIN Vt ;
: DRAW-PTS (hi lo -) DO I PTX DRAW LOOP ;

The above code is essentially all there is to it. The remainder of
the screens illustrate using the drawer routines.

---

Screen # 3

: USAFRAME
(draws a double border around the map area)
GRAY FG 320 42 AT 631 42 DRAW 631 196 DRAW 320 196 DRAW 320
42 DRAW WHITE FG 316 38 AT 635 38 DRAW 635 200 DRAW 316
200 DRAW 316 38 DRAW ;
: FILLUSA 480 100 GREEN GRAY FLOOD ;

---

Screen # 4

: REGION# 4112 GOTOXY GRAY FG
(prints # of regions)

"Regions of approx. similar rainfall." WHITE FG
ASCII 1 130 80 ORIGIN V+ GENIT
ASCII 2 110 45 ORIGIN V+ GENIT
ASCII 3 144 36 ORIGIN V+ GENIT
ASCII 4 156 6 ORIGIN V+ GENIT
ASCII 5 68 14 ORIGIN V+ GENIT
ASCII 6 56 50 ORIGIN V+ GENIT
ASCII 7 25 23 ORIGIN V+ GENIT ;

Note: GENIT is PC/FORTH for setting a character in graphics
mode at x y coordinates rather than column/row locations.

---

Screen # 5

/this sequence will open/close required files.

: DRAWMAP
(to draw the USA...)
: USAFRAME
(draw frame around image area)
: CLOSE-SCR
(close current screen "drawer")
: EMPTY-BUFFERS
(clean out any blocks stored)
: POINTFILE OPEN-SCR
(open the screen-file "pointfile")
: GREEN FG DRAMUSA FILLUSA
(draw the stuff...)
: RED FG DRAMSTEEL .REGIONS
: CLOSE-SCR
(close the drawing-file)
: EMPTY-BUFFERS
(clean any junk from buffers)
: DRAWER OPEN-SCR
(re-open "drawer")
2DUP ;

Note: GENIT is PC/FORTH for setting a character in graphics
mode at x y coordinates rather than column/row locations.
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WITH THE CONVENIENCE
OF CONVENTIONAL LANGUAGES

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Screen # 6

( DEMO: Setup jbb 14:52 04/20/86 )
: DEMO: Setup jbb 14:52 04/20/86 )
: HEADER 0 0 GOTOXY ; microARCH Drainage Module ;
: GIVEN ; Given the following information ;
: SETUP 840X350 VMODE ; (EGA high-res graphics mode)
BLUE 86 CLS
WHITE FG HEADER 0 14 639 14 LINE 0 233 639 233 LINE
RED FG 0 15 639 15 LINE 0 232 639 232 LINE
GRAY FG 0 16 639 16 LINE 0 231 639 231 LINE ;

Screen # 7

( Dummy variables for demo jbb 11:18 04/29/86 )
2VARIABLE D 23400 0 D 2!
2VARIABLE SLOPE 340 0 SLOPE 2!
2VARIABLE AC 125 0 AC 2!
2VARIABLE C 45 0 C 2!
VARIABLE REGION 3 REGION !
VARIABLE YEAR 10 YEAR !

\ this is a dummy floating point F and formatted F for DEMO
: F ? 20 <! ASC ! ASCII . HOLD 85 > TYPE SPACE ;

>ENTRY-W 0 17 GOTOXY ; ( dummy window control for DEMO )

Screen # 8

( Problem jbb 15:21 04/20/86 )
: DEMO SETUP DRAWMAP WHITE FG 0 3 GOTOXY GIVEN GRAY FG
1 5 GOTOXY ; Distance of flow (ft.) ; * D F?
1 6 GOTOXY ; Slope ( %) ; * SLOPE F?
1 7 GOTOXY ; Coefficient of runoff ; * C F?
1 8 GOTOXY ; Region (Steel Chart) ; * REGION ?
1 9 GOTOXY ; Design Storm (year) ; * YEAR ?
1 10 GOTOXY ; Area (acres) ; * AC F?
>ENTRY-W WHITE FG
" What is the Volume of runoff (Q) in cfs ? " GRAY FG CR ;

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Synonyms and Macros, Part 4

Compiler Macros

Victor H. Yngve
Chicago, Illinois

The important thing about compiler macros is that they postpone the normal compile-time action of words to a later compile time, thus providing us with valuable programming options. Let me explain.

The macro definitions written

MACRO <name> . . . END-MACRO

that were previously introduced may be called compile-and-CMOVE macros. They can be used to increase readability where a colon definition cannot be used because of interference with the return stack, e.g., in accessing loop indices, and they can be used to increase execution speed in time-critical parts of the code while preserving the readability of colon definitions.

However, there is the possibility of further improvements in readability where compile-and-CMOVE macros cannot be used because they would contain unpaired compiler words like IF, BEGIN or LOOP, or the Forth implementation has absolute rather than relative branch addresses. These cases would have to be coded in the familiar manner by what we can call COMPIL-[-COMPILE] compiler macros: Each nonimmediate word is preceded by COMPILE, each immediate word by [COMPILE] and IMMEDIATE is appended to the definition. Some examples are given on screen fifty-eight, which shows macros for IF-NOT and WHILE-NOT suggested by Ed Petsche.

But COMPIL-[-COMPILE] compiler macros can be confusing and error prone. One has to remember which words are immediate and which not in order to choose between [COMPILE] and COMPILE, and it is easy to forget to append IMMEDIATE. Also, the resulting definitions are not very readable, thus tending to defeat an important reason for their use.

COMPIL-[-COMPILE] compiler macros are even more difficult and awkward to use if the definition contains numbers. These were not covered in Jeffrey Soreff's article, which warrants close study, and it may not be immediately clear how to handle them. The use of numbers in COMPIL-[-COMPILE] compiler macros is shown on the same screen in a suggested macro IF# for testing whether the number on the stack represents the ASCII code for a decimal digit.

The macro facility introduced here can take care of these cases by the words << (postpone) and >> (end postpone). Simply write

: <name> . . . << . . . >> . . . ;

with the words of the macro enclosed between the angle brackets. It is not necessary to append IMMEDIATE to the definition. A Forth-83 style glossary entry is provided on screen sixty-two. These << . . . >> macros are called angle-bracket compiler macros. They do nothing more than can be done with COMPIL-[-COMPILE] compiler macros, namely they make a compiler word that postpones the immediate or nonimmediate action of the enclosed words, but they do it automatically. Their use is illustrated on screen sixty-three in the recoding of IF-NOT, WHILE-NOT and IF# as angle-bracket compiler macros. The increased readability due to removing the COMPILE and [COMPILE] clutter is especially apparent in the case of longer definitions like IF#.

Angle-bracket compiler macros can be nested in normal fashion together with colon definitions and compile-and-CMOVE macros. However, one cannot nest

[ . . . ]

inside an angle-bracket compiler macro to postpone the run-time action of the enclosed words. Use

>> . . . <<

instead, thus dividing the compiler macro into two and compiling the included words normally. Note that

: <name> . . . << . . . >> ;

is a postponing version of

The main purpose of using compiler macros is to postpone the execution of immediate words when this is necessary for resolving branch addresses or to postpone the processing of the input stream. Their main disadvantages for other purposes over compile-and-CMOVE macros are that they take up more space in the dictionary and they cannot be tested directly at the console without first embedding them in other words.

In a compiler macro, immediate words have their normal, immediate compile-time action postponed to when the compiler macro is used in compiling another word. With words like ." this means compiling a string from the input stream for later output. With words like IF and LOOP this means compiling branches and resolving branch addresses. With IF# this means compiling the compilation addresses of a word found in the input stream. For LITERAL it means compiling a number on the stack as a literal. In all these cases the normal run-time action resulting from the compile-time action is postponed to the time when the other word is executed. But for immediate words like + and ." nothing is compiled at the postponed compile time, for their normal compile-time action is to be executed.

Nonimmediate words such as SWAP, +, CREATE, .(tick) and .(comma) have their compilation postponed to the time when the compiler macro is used in compiling another word. Their run-time action is then postponed to the time when this other word is executed. At that time, .(tick) finds a Forth word in the input stream and leaves its compilation address on the stack, and words like CREATE make the following text in the input stream into the name field of a word being defined.

Thus there are three times associated with compiler macros where information is processed into a different form: a postpone time when each immediate word is compiled and each nonimmediate word is compiled with a preceding COMPILE, a compile time when the normal immediate or nonimmediate
action of the words is taken in compiling another word, and a run time when the other word is executed. In this they are like SYNONYM, which has a time when the synonym is defined, a time when the synonym compiles the original word and a time when that word is executed.4

And in this they are different from compile-and-CMOVE macros, which have only a compile time and a run time, the run time being delayed by simply moving the compiled words into place with CMOVE without further processing. There are also input-stream macros, which have been explored by Don Taylor.5 They also have only a compile time and a run time, but here the compile time is delayed by saving a copy of the source text in the dictionary.

In brief, the implementation of << and >> is straightforward and can be done entirely in Forth-83 with the help of the familiar nonstandard words defined in Forth-83 on screen fifty-nine. The word << (postpone) simply marks the word in which it appears as immediate, looks up each following word in the dictionary, determining whether it is immediate or not, and compiles it, preceding it by COMPILE if it is not immediate. This gives the same result as using a COMPILE-[COMPILE] compiler macro. If a word is not found, the auxiliary word <NUMBER> tries to convert it to a number and compile it appropriately as a literal with a following [COMPILE] LITERAL. Compilation is terminated by detecting >> (end postpone) in the input stream. The stretch of code between << and >> is marked to be used only while compiling.

For ease of learning and ease of use, it is important that numbers be treated the same as by the outer interpreter, which usually handles both double- and single-precision numbers. The standard does not dictate how double numbers should be represented in the input stream. The system implemented here in <NUMBER> is that if the number contains a period either initially, finally or in the middle, it is compiled as a double number. No record of the location of the decimal point is kept. Ideally <NUMBER> should be modified to use the double-

Screen # 56
0 ( Examples of macros using COMPILE and [COMPILE] )
1
2 ( from Ed Petsche, Forth Dimensions VII,3:6 )
3 : IF-NOT COMPILE 0= [COMPILE] IF ; IMMEDIATE
4 : WHILE-NOT COMPILE 0= [COMPILE] WHILE ; IMMEDIATE
5
6 : IF# COMPILE DUP #47 [COMPILE] LITERAL COMPILE >
7 COMPILE OVER 56 [COMPILE] LITERAL COMPILE <
8 COMPILE AND [COMPILE] IF ; IMMEDIATE
9
Screen # 59
0 ( Extensions needed to Forth-83 Required Word Set )
1 32 CONSTANT BL
2 : ASCII (-- n ) BL WORD 1+ C8
3 STATE # IF [COMPILE] LITERAL THEN ; IMMEDIATE
4 : ?COMP STATE # 0= ABORT" Compilation only "
5 : ?PAIRS - ABORT" Conditionals not paired "
6
Screen # 60
0 ( Compiler Macros <NUMBER> )
1 vhy 11/14/85
2
3 2 : <NUMBER> ( 0 0 addr -- ) ( Postpone Number Or Abort )
4 DUP 1+ C8 ASCII = - DUP >R ( neg flag to return stack )
5 IF 1+ THEN
6 CONVERT DUP C@ ASCII = DUP >R ( doub ) ( lo hi ad : df nf )
7 IF CONVERT THEN C@ BL = ( found? ) ( lo hi ff : df nf )
8 IF R> R> SWAP >R ( neg? ) ( lo hi nf : df )
9 IF DNEGATE THEN ( lo hi : df )
10 SWAP [COMPILE] LITERAL COMPILE [COMPILE] LITERAL R>
11 IF [COMPILE] LITERAL [COMPILE] LITERAL [COMPILE] LITERAL
12 ELSE DROP
13 THEN
14 ELSE 1 ABORT" Not found "
15 THEN
16
Screen # 61
0 ( Compiler Macros << ... >> )
1 vhy 11/14/85
2
3 1 : >> ( -- ) 1 ABORT" Unpaired " ; IMMEDIATE ( End Postpone )
4 2 : << ( -- ) ( Postpone )
5 ?COMP
6 COMPILE ?COMP ( this word compile only )
7 IMMEDIATE ( same for macro words )
8 BEGIN
9 BL WORD FIND ( compiled word immediate )
10 END
11 IF COMPILE THEN ( search dictionary )
12 OVER [ | ] >> = NOT ( not done? )
13 IMMEDIATE ( was word found? )
14 IF 1- ( was it nonimmediate? )
15 IMMEDIATE ( for nonimmediate words only )
16 ELSE DROP <NUMBER> ( postpone number or abort )
17 RESTORE DROP DROP ; IMMEDIATE

Screen # 62
0 ( Compiler Macros Glossary entry )
1 vhy 11/14/85
2 "postpone"
3 A compiler word used in the form:
4 : <name> ... < ... > ... ;
5 When << is used while compiling, it renders the definition
6 immediate, compiles each immediate word between << and >> ,
7 compiles each nonimmediate word with a preceding COMPILE,
8 compiles each single number as a literal followed by
9 [COMPILE] LITERAL, and each double number as two literals
10 each followed by [COMPILE] LITERAL. The stretch of words
11 between << and >> are marked compile only. The result is
12 that the immediate or nonimmediate action of the words and
13 numbers between << and >> is postponed by one step of compila-
14 tion. The word >> (end postpone) is freestanding and
15 surrounded by spaces.
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Figure One

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<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Figure Four

Volume VIII, No. 3 27 FORTH Dimensions
number scheme of the implementation in which it is installed.

Angle-bracket compiler macros have a variety of programming applications. For example, they are convenient for building new programming constructs. Their use in defining the looping words $\text{FOR}$ and $\text{NEXT}$ is illustrated on screen sixty-three. These words repeat a series of words $n$ times by writing

$\text{FOR} \ldots \text{NEXT}$

The index $n$ is counted down to zero, changes the value of $n$ and is available by using it inside the loop.

On the same screen, a high-level definition of the word $\text{TO}$ (sometimes called $\rightarrow$ or $\mathbf{is}$) is given using a compiler macro. This word works with special constants defined by $\text{VALUE}$ (given on the same screen) which can have their value changed. One writes

$\text{VALUE \SCORE}$

Then, executing

$\text{16 \TO \SCORE}$

changes the value of $\text{SCORE}$ to sixteen. When the word $\text{TO}$ is used, it obtains the address of the body of $\text{SCORE}$ with

Screen # 65

$0 \text{ ( An implementation of Eaker's CASE: Forth Dimensions II/3:37 )}$

$1$ $2 \text{ VALUE OLD-DEPTH ( Eaker used the fig-FORTH CSP for this. )}$

$3$ $4 \text{ CASE } ?\text{COMP OLD-DEPTH DEPTH TO OLD-DEPTH } 4 \text{ ; IMMEDIATE}$

$5$ $6 \text{ OF } 4 \text{ ?PAIRS } \text{ << OVER } = \text{ IF DROP } >> 5 \text{ ;}$

$7$ $8 \text{ END} \text{ OF } 5 \text{ ?PAIRS } \text{ << ELSE } >> 4 \text{ ;}$

$9$ $10 \text{ ENDCASE 4 ?PAIRS } \text{ << DROP } >>$

$11$ $12 \text{ BEGIN DEPTH OLD-DEPTH } \not = \text{ WHILE } \text{ << THEN } \text{ >> REPEAT}$

$13$ $14 \text{ TO OLD-DEPTH } ; \text{ ( restore depth-value for nesting )}$

$15$ $16 \text{ ENDCASE CR }$

Screen # 66

$0 \text{ ( Macro Sieve 1 Preface Version E whv 12/5/85 )}$

$1$ $2 1910 \text{ CONSTANT SIZE CREATE FLAGS SIZE ALLOT}$

$3$ $4 \text{ MACRO SET-FLAGS-TRUE}$

$5$ $6 \text{ FLAGS SIZE } 1 \text{ FILL END-MACRO}$

$7$ $8 \text{ SYNONYM OCCOUNT}$

$9$ $10 \text{ DO-FLAGS } \text{ << SIZE } 0 \text{ DO } >> ;$

$11$ $12 \text{ IF-TRUE } \text{ << FLAGS I } + \text{ C0 IF >> ;}$

$13$ $14 \text{ MACRO GET-PRIME } ( - p ) \text{ I DUP } + 3 + \text{ END-MACRO}$

$15$ $16 \text{ MACRO FIRST-MULTIPLE } ( p - p m ) \text{ DUP I } + \text{ END-MACRO}$

$17$ $18 \text{ WHILE } < \text{ SIZE } ( m - m ) \text{ << DUP SIZE } < \text{ WHILE } >> ;$

$19$ $20 \text{ MACRO SET-FALSE } ( m - m ) \text{ O OVER FLAGS } + \text{ C1 END-MACRO}$

$21$ $22 \text{ MACRO NEXT-MULTIPLE } ( p m - p n ) \text{ OVER } + \text{ END-MACRO}$

$23$ $24 \text{ SYNONYM DROP-MULTIPLE DROP}$

$25$ $26 \text{ MACRO DROP-MULTIPLE DROP}$

$27$ $28 \text{ MACRO INC-COUNT}$

$29$ $30 \text{ NEXT-FLAG } \text{ << THEN LOOP } >> ;$

$31$ $32 \text{ PRINT-COUNT" Primes " } ;$

Screen # 67

$0 \text{ ( Macro Sieve 2 Algorithm Version E whv 12/5/85 )}$

$1$ $2 \text{ MACRO CANCEL-MULTIPLES } ( - p \text{ -- prime )}$

$3$ $4 \text{ FIRST-MULTIPLE}$

$5$ $6 \text{ BEGIN WHILE } < \text{ SIZE SET-FALSE NEXT-MULTIPLE REPEAT}$

$7$ $8 \text{ DROP-MULTIPLE END-MACRO}$

$9$ $10 \text{ DO-PRIME}$

$11$ $12 \text{ SET-FLAGS-TRUE OCCOUNT}$

$13$ $14 \text{ DO-FLAGS}$

$15$ $16 \text{ IF-TRUE GET-PRIME CANCEL-MULTIPLES PRINT-PRIME INC-COUNT}$

$16$ $17 \text{ NEXT-FLAG}$

$18$ $19 \text{ PRINT-COUNT" Primes " } ;$

$20$ $21 \text{ FREE FORTH CATALOG}$

$22$ $23 \text{ MOUNTAIN VIEW PRESS}$

$24$ $25 \text{ PO BOX 4656}$

$26$ $27 \text{ Mountain View, CA 94040}$
>BODY and, if executing, uses it to store the number on the stack into the value. But if compiling, it instead compiles the address as a literal and then compiles CSTORE, thus postponing the actual change of the value to the word of the time in which TO SCORE is contained. Thus a new value on the stack can be stored into SCORE with the speed of execution of

\[ \text{addr} \]

where \text{addr} is compiled as a literal.

An exercise in dealing separately with postpone time, compile time and run time is given on screen sixty-four. Suppose there is an event in which new world records are continually being made. The variable \text{RECORD} is updated at each of these times by storing in it the current new world record. If one then executes

\text{n AWARDS?}

at run time, the routine will successively test \text{n} against the different records stored at postpone time, at compile time and at run time, giving a printout like

\text{FOR, NEXT, TO} and \text{CASE} constructs, for these could be defined once and for all by a Forth system programmer as \text{COMPILE-[COMPILE]} macros. Their poor readability would be of little concern. But the genius of Forth is that everyone can be his own system programmer. The essence of Forth programming for anything but the simplest program is to invent an appropriate language tailored to the application, program the application in it and implement the special language in terms of Forth words in a prologue.

Thus the most important uses of angle-bracket compiler macros is in everyday programming to assist this process of organizing a program into appropriately named pieces. As an illustration, screens sixty-six and sixty-seven show the sieve benchmark program rewritten from version D in the last article, here using angle-bracket compiler macros as well as compiler-CMOVE macros, synonyms and colon definitions. The further increase in readability is due to using angle-bracket compiler macros to name coherent stretches of code that happen to include unpaired conditional and looping words. If the implementation has absolute rather than relative addresses, the word \text{CANCEL-MULTIPLES} should be written as an angle-bracket macro as well. There is no increase in the running time of the benchmark, for exactly the same code is compiled as by the originally published sieve algorithm.

Let us conclude with a summary of when these various different types of definition should be used in the course of general Forth programming. Use:

\text{<name> . . . ;}

for all normal purposes in Forth with the following exceptions:

Use

\text{SYNONYM <name> <name>}

for any definition with only one word in it. Synonyms run faster, take up less room in the dictionary and can be used for any word, whether it is immediate or not.
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If: <name> ... ; cannot be used because the return will interfere with the return stack, or if it is desired to save the time of nesting and unnesting at run time, use

MACRO <name> ... END-MACRO

with the following exceptions:

If: <name> ... ; or MACRO <name> ... END-MACRO cannot be used because the code contains unmatched compiler words like IF or DO, or if MACRO <name> ... END-MACRO cannot be used because it contains branches compiled by words like IF or LOOP and the Forth implementation being used has absolute rather than relative branches, use

: <name> << ... >> ;

instead.

In addition, use

<< ... >>
in colon definitions along with other words to make compiler words that postpone compilation of the enclosed words for special purposes.

References
Shuffled Random Numbers

Leonard Zettel
Trenton, Michigan

In my opinion, the random number generator given by Leo Brodie in *Starting Forth* is indeed a very good one. While I have not checked it exhaustive-ly, it seems to follow all the recommendations in Knuth\(^1\), which is usually regarded as the definitive work on random number generation. It could therefore be regarded as best in its class, the class being sixteen-bit linear, congruential, pseudo-random number generators (if you really need to know what that means, read the reference).

That said, it must also be said that it shares a weakness shown by all generators of its type. To highlight the problem most dramatically, use `CHOOSE` to generate \(x,y\) pairs and plot them on your computer screen. When I do this on my Commodore 64 (using Forth extensions that come with SuperForth 64 by Parsec Research), with \(x\) going from 0 to 319 and \(y\) going from 0 to 159, what I get looks fine at first but eventually settles down to a tight corduroy pattern of diagonal lines. Roughly half the points never get plotted at all!

Now, if I ever see raindrops doing that on the pavement, I would get highly suspicious. Random behavior it ain't, though possibly good enough for many purposes to which random numbers are applied. When I used the numbers to study random walks, they turned out to be not good enough, because the walk kept insisting on avoiding large portions of the screen. In general, you will run into trouble if you use the numbers in bunches of \(n\), with the first, second, etc. always used for the same purpose. In the plotting example, there are bunches of two, with the first number always used as the \(x\) coordinate and the second as the \(y\). In general, you can expect the problem to get worse for larger \(n\).

Fortunately, Forthwrights, there is a solution, which Knuth gives on page thirty-two. The answer is to shuffle the numbers. The screen shows a Forth implementation of his algorithm B. We put thirty-three random numbers generated using `RANDOM` into the array `RANDARRAY`

using `RANDARRAY.INITIALIZE`. Then, when we need a number we invoke `SHRANDOM` to get a shuffled, random number. We use the thirty-third number in the `RANDARRAY` to calculate a position in the array. We then fetch to the stack the random number at that position, replacing it in the array with a new number from `RANDOM`. We then `DUP` the random draw we have on the stack and store one copy at the end of `RANDARRAY`, leaving the other on the stack as the result of `SHRANDOM`.

For good measure we have `CHOOSE`, which works just like Brodie's `CHOOSE` except that it uses `SHRANDOM` instead of `RANDOM`. For those of you tempted to tinker and, maybe, tempted to replace the multiply with a shift operation of some kind, note that `CHOOSE` and `SHRANDOM` use the leftmost bits of the random number to generate the desired integer. This is recommended by Knuth, since the leftmost bits of the output of a congruential generator are, in many senses, "more random" than the right-hand bits.

Now for some comments and, if I can get away with it, some editorializing. I do not personally know of any way to get better quality random numbers (with a sixteen-bit modulus) than the screen presented here. On the other hand, being a somewhat lazy amateur, I have not tested the output nearly as thoroughly as I might have. In fact, the only tests I have run are the plots and the random walk. In addition, it is worth quoting Knuth (page 173), "Perhaps further research will show that even the random number generators recommended here are unsatisfactory; we hope this is not the case, but the history of the subject warns us to be cautious."

If Forth is indeed the best computer language in existence, and I firmly believe that it is, then it deserves to have only the very best routines put forward as part of its repertoire. You readers out there are the Forth community, and collectively you have more knowledge and experience than any single author. It would not surprise me at all if someone out there could improve on what I have presented. If you can, you have a duty to let the rest of us know about it; that is what communities are all about. If you don't want to work it up for publication, I would be glad to hear from you personally.

Reference

The Multi-Dimensions of Forth

Glen B. Haydon
La Honda, California

Forth intrigued me when I first came across it six years ago. I had about ten years experience programming in a variety of high-level languages and some exposure to assembly languages. It always seemed that I needed to do things that were not immediately possible with the tools available in those languages. Furthermore, many of the languages required a mainframe computer for the compiler alone.

An acquaintance I met in a computer store dragged me along to a fourth Saturday meeting of the Forth Interest Group. I caught the disease. Many others have caught the disease. The symptoms vary, but its contagious nature is there.

I have tried to understand just what it is about Forth that is so contagious. I have found that Forth has many dimensions. Forth is a religion. Forth is a philosophy. Forth is a software kernel which emulates a hardware design.

Forth is a hardware processor. Forth is the assembly language for that processor. Forth is an operating system. Forth is a high-level language. Not all of these dimensions intersect.

Forth as a Religion

Franz Werfel in the introduction to the “Song of Bernadette” has the following quotation:

“For those who believe in God, no explanation is necessary; for those who don’t, no explanation is possible.”

Those who believe in Forth have no problem with the requirements of reverse Polish notation, stack manipulations and the many unconventional ways of Forth. There is no way to explain Forth to experienced programmers who find the unusual characteristics of Forth impediments to “good” programming. A huge gap exists between the two points of view.

The newcomer to Forth is often better off if he has had no previous computer experience. After some time, he either becomes a member of the sect or he rejects it and goes on to something else. Those who end up using Forth let it become a part of their life.

Unfortunately, the religious image of Forth has been played up by the public media. It has had a way of turning off those who already have a religion.

Forth as a Philosophy

I find it difficult to pronounce “!”. It is certainly not a common phonetic symbol. But “store” is a common Forth function. The symbol is like Chinese ideograms with no phonetic meaning.

I found that Forth functions are like the concepts associated with Chinese ideograms. I wrote a paper for the 1981 Rochester Forth Conference, "Forth and the Nature of Ideographic Thought.” There is something very different between the way the Forth ideograms are assembled to convey an idea and more conventional ways of the classical high-level programming languages. The root of programming in Forth is to rethink the problem. Don’t use thought processes derived from experience with the other programming languages.

After working closely with Charles Moore for about six months, I observed another dimension of Forth. Charles Moore has a way of immersing himself in the problem at hand. There was no such thing as an eight-hour work day. His entire life was immersed in the problem for days at a time.

Of course, his problems included the use of computers. Charles Moore attempted to understand all aspects of the problems. Many have heard him talk about the unreliability of hardware. He had to learn how to make the hardware, if possible. At the other end of programming computer applications, he found that most end users did not really understand the problem they were trying to solve. His approach was to immerse himself in all aspects. There were occasional escapes into science fiction and dreams of what things might be.

This total immersion in problem solving seems analogous to the ways of Zen. From time to time I have added pieces to a notebook which I have labeled, “Zen and the Art of Problem Solving.” The title alone captures the spirit of the Forth philosophy.

Forth as a Software Kernel

The origins of Forth were based on trying to make existing hardware do the things that the applications required. It was a sort of virtual machine.

The Forth kernel simulates a small, necessary and sufficient set of functions which allows one to get on with the problems at hand. At the center of Forth implementations is some sort of kernel which allows the user to do what he needs to do. That Forth has been implemented on such a wide variety of processors is testimony to the virtue of a small, necessary and sufficient set of functions. It is a relatively simple matter to move the kernel from hardware system to hardware system.

A careful review of the necessary and sufficient functions shows that somewhere between sixty and seventy functions are convenient. It has been said that only eleven or twelve primitives programmed for a specific processor will provide all of the necessary operations with which to build a full, convenient set. But most implementations of Forth kernels include thirty or more specific functions unique to the specific hardware.

A unique design feature of Forth is its use of multiple stacks. Multiple stack architectures have been discussed since the mid-1950’s. No true hardware implementation of a multiple-stack architecture was produced. Burroughs did provide an optimized set of pointer registers to utilize a part of main memory as an efficient stack. But the dedication of high-speed RAM to each stack apart from main memory is a relatively recent turn of things.

Forth as a Hardware Processor

There is a better way than emulating Forth in hardware: design the hard-
ware to do the job. Many of the problems of implementing Forth on other processors just go away. Furthermore, a dramatic speed increase is possible — not just a factor of two or four, but a factor of one or two orders of magnitude.

Hartronics has produced a bit-slice design. Their product is proprietary but very efficient. In the last year or so, several other groups have addressed the problem of implementing a Forth kernel in hardware; not the least of these efforts has been that of Charles Moore. His group designed an architecture and committed it to silicon as the Novix 4000 chip. Several vendors are including that chip in systems. The systems are generally small and have low power requirements.

A group at Johns Hopkins have taken the basic design of the Novix 4000 device and expanded it to a thirty-two-bit processor on a chip. Again the functions are hard wired into the silicon. A group in Hull, England, formed a company named Metaforth. They have produced a very fast board implementation which is becoming available. It has been optimized with custom components and is proprietary.

I have been associated with yet another design of hardware. The design features a writable control store and dedicated hardware stacks. The writable control store can be used to implement Forth or to experiment with other programming languages. It is able to run a functional Forth system as fast as those using a chip. The product is available today as a kit. All details are included in the wire-wrap kit to allow assembly from inexpensive, garden-variety components.

A writable control store machine is a true extension of the Forth philosophy of extensibility. It makes it possible to redefine and extend the actual processor functions. One should think of these machines as another processor design. The machine language for that processor can be made to execute a Forth kernel.

Forth as an Assembler

In conventional thinking, a Forth assembler provides access to the processor on which a software kernel is implemented. When the processor executes Forth functions, Forth becomes the assembler mnemonics for the processor. There simply is no other assembler.

With this level of maturity, the functions of the Forth kernel are directly compiled to the machine language of the Forth processor. As with many other processors, various fields in the machine language are directly coupled with the function of the hardware components.

Thus, one dimension of Forth is comparable with the assembly languages used with other processors.

Forth as an Operating System

Charles Moore found all of the overhead of an underlying operating system with its file structures and so on, completely unnecessary for his applications. He simply did not use the operating system’s facilities in his work. Other early Forth users felt much the same way. They learned to work with their rudimentary operating system and found it most satisfactory. Forth was the operating system and the entire external storage was randomly accessed as a single file. What could be simpler?

But then there came those experienced with other operating systems and other programming languages. They expected to have available certain facilities. And many of them were using their hardware for other applications as well. They needed to have things “compatible.”

The hardware emulation or simulation of a multiple stack system in Forth is not sufficient to yield a fully functional system.

When I started with the kernel from the *fig-FORTH Installation Manual*, everything seemed to work fine. I could load screens but I could not write to screens. Only if I were able to get something on a screen which was compatible with my drives, could I load it.

Interactively, I had to type in the necessary functions to allow text to be placed on a screen. Or, I could write the necessary functions in assembly language and have them available in my

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FORTH as a High-Level Language

Before I caught the Forth disease, I had a moderate amount of experience with a variety of high-level languages on mainframe computers. I found that PL/I provided me the most access to the functions I needed. But the PL/I compiler was large! There was no way it could be reduced to the 1K system I started with. Gradually, the small systems got larger, but a full implementation of PL/I would not even fit on many small machines today. Furthermore, even PL/I would not do everything I thought I wanted to do.

As a high-level language, the fig-FORTH Model left much to be desired. Many of the functions I expected just were not there. String functions were missing. There was no floating-point package. But I found that a fig-FORTH Model with a few custom bells and whistles served me well.

Then came a change of functional definitions for forty of the words in the fig-FORTH Model (suggested by the Standards Team). I really did not see much advantage to many of the changes. There was one major one in my book — the change of the CREATE ... DOES> function.

Reluctantly, I went ahead and made the forty changes to conform. This had the advantage that there were soon two tutorials — Leo Brodie’s Starting Forth and Alan Winfield’s The Complete Forth. The revision of the fig-FORTH Model with the forty functional changes became MVP-FORTH. I summarized my learning notes in a glossary that I am glad others have found useful.

At that point, I had a fully functional Forth which allowed me to go on with application programming. MVP-FORTH has remained stable for over four years. A substantial library of Forth functions is now available which run on the model I use. I will not be confined by the additional stipulation imposed by the 79 Standard. MVP-FORTH is not 79 Standard. I have no time to make the changes suggested by the 83 Standard and will not be confined by the additional stipulations.
My experience is not unique. Many other vendors have developed excellent implementations of a functional Forth. Some of them have kept their kernel and parts of the rest of their systems proprietary. This is fine. They have provided excellent documentation for their products. This is a big help. However, most other vendors have made regular revisions. Some of the books on Forth for some vendor's dialects are no longer correct.

Now there is little point in changing your functional Forth just to be changing. Of course, there are many implementation techniques which could be adopted. One can use direct-threaded code or indirect-threaded code or token-threaded code. They make little difference in the application of Forth to specific projects. There is little point quibbling about the details. There are better things to do.

After learning Forth, one should adopt a kernel and add only the necessary and sufficient routines to provide a uniquely functional Forth. For example, add your favorite editor and, only if you need it, add the floating-point functions. Collect your own library of Forth “goodies.” The ability to do this is what is attractive about Forth.

As a high-level language, Forth is as extensible as you wish to make it. However, you need not be burdened with unnecessary and unused extensions.

Conclusions

I have taken you on a short tour of my attempts to understand Forth and where it is going.

Whatever you may want to call the process of immersing yourself in all aspects of a problem — a religion, a philosophy, even Zen — Forth embodies an approach to problem solving.

The essence of the Forth language is its kernel — an emulation of a hardware machine. As such, it provides a new approach to processor design. The new hardware machines with multiple stacks will open new doors to the computing world. They will provide new alternatives to the conventional processors and the relatively new RISC machines.

If the primitive nature of a simple Forth operating system is satisfactory for your application, there is no reason to add the complexities of other systems. If not, you are free to add whatever operating system you wish.

Forth as a high-level language will shake down and take its place along with everything from BASIC to ADA. As has been demonstrated, one can even implement other high-level languages on a Forth software or hardware kernel. PROLOG has been implemented in a LISP which was implemented in Forth. I don’t know what to call that combination.

The dimensions of Forth encompass many different aspects for any particular user. It is a little like Humpty-Dumpty in Alice in Wonderland. You need to make clear which combination of dimensions you are talking about, and to understand which others are talking about.

It is the richness of these many dimensions that makes Forth what it is.

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Stack Numbers by Name

Melvin Rosenfeld
Santa Barbara, California

When a word involves more than two or three numbers on the stack, the logistics of accessing them is often tedious. Sometimes the natural idea of putting the numbers in variables is not suitable. An example of this will be given shortly (example two).

I would like to offer a solution that I’ve devised. It consists of two words: \{ (the left brace) and \; (a double semicolon). Below are two illustrations of how these words are used.

Example One A definition of the function

\[ f(x,y,z) = (x^2 + y)(xz + y) \]

is shown in Figure One. In that figure, the three words X, Y and Z that appear between the braces can be arbitrary. They are dummy names for the relevant stack entries. They will be in the dictionary only during the compilation of the word F. The stack effect of F will be \( (X \ Y \ Z -- \text{-- answer}) \).

Example Two As a consequence of the method of implementing the use of names for numbers on the stack, the function being defined can be used recursively. In this example, a word which solves the “Towers of Hanoi” puzzle will be given. I won’t bother to remind readers what the puzzle is, since the point to be made is how neatly a recursive word can be defined. In this word, it is presumed that a word OUTPUT has already been defined (see Figure Two). OUTPUT takes two numbers a and b off the stack and reports, say by printing on the screen, “Take the top disk off peg a and place it on peg b.”

In this example, one cannot use a variable to hold “the” value of, say, \#DISKsin since many such values must be kept while the recursive function HANOI wends its way down.

Discussion of the Method

Since \{ and \; are used to define a word, there are two aspects to consider: first, the run-time effect of their use on the word being defined, and second, their operation during compilation of the word. Let us consider the word F from example one. Prior to its compilation, it is necessary to have the following words in the dictionary:

```
0 VARIABLE STACKTOP
: MARKSTACK SP@ STACKTOP @;  
: C1 STACKTOP @ 0 + @;  
: C2 STACKTOP @ 2 + @;  
: C3 STACKTOP @ 4 + @;  
```

The dictionary entry for F (after it has been compiled) will, in its parameter field, contain the sequence of CFA’s of the words comprising its definition. The names of this sequence of words are, in order,

```
MARKSTACK C3 DUP * C2 + C3 C1 * C2 + *  
```

then a literal, in this case three (3), followed by the CFA of another word, SHIFTSTACK and finally \;s.

The word SHIFTSTACK must also be in the dictionary before F can be compiled. Its name suggests what it does, but its definition is best left until later. Since the temporary name X for the third number on the stack has been replaced by C3, there is no reason for X to be in the dictionary.

Next we consider the compilation phase of the word F. Consider the string of symbols

```
: F { X Y Z } \;s
```

When the interpreter reaches : it of course begins to compile the next word, namely F. It is thus in compile mode when \{ is encountered. \{ is an immediate word and, instead of being compiled, does the following:

1. Concludes the definition of F (so that, for the moment, it is a no-op).
2. Counts the number of words in the input stream until \} (three words, in this case).
3. Defines the first word, namely X, to be an immediate word. The action of X will be to place the CFA of C3 into the next place in the dictionary. It similarly defines Y and Z.

4. Begins compilation of the next word in the input stream, namely \}. The interpreter continues to compile the words in the input stream until \; is encountered. This latter word is also immediate. It does the following:

```
1. Compiles the literal 3 and one last word, namely SHIFTSTACK into the definition of \}.
2. Concludes the definition of \}.
3. Shifts within the dictionary the parameter field of \} so that it abuts the parameter field of F with the effect that the definition of \} becomes the definition of F.
4. Erases from the dictionary the words X, Y, Z and \}.
```

There is one additional word that may be useful, namely \;s. It acts very much like \; but doesn’t cause the stack to be shifted at the conclusion of the run-time action of F.

Time Penalty for Naming Numbers

The time spent in compiling seems unimportant, since it is undoubtedly made up for in time saved during programming. The run-time cost is important (if the word will be used many times in a program). Figure Three shows a short test to check the overhead of using this technique.

The overhead — the difference between right way and wrong way — is 43/20000 seconds per implementation, about 1/465. Most of the time is taken up by the word SHIFTSTACK. Figure Four provides a second test, in which \;s was used and, therefore, SHIFTSTACK isn’t needed. The overhead per implementation without SHIFTSTACK is 1/1176 seconds.

Since only a few words (\;, MARKSTACK, SHIFTSTACK, C1, C2, C3 ... C10) are used during run time, if they were written in machine language, I suspect test2 would hardly be slower than test1 and there would be little penalty for using what seems to me a very natural way of accessing the stack.
The version of Forth that I’ve used is Wyco FORTH for the TI 99/4. I believe it is a variant of fig-FORTH. I have used the word move which in this Forth seems to diverge from the standard.

In Wyco FORTH:

```forth
addr1 addr2 count MOVE
```

Moves the block of main memory bytes from addr1 to addr2. The length of the block in bytes is given as count. The direction of the move is determined so that no overlap occurs, with the word CMOVE or <CMOVE called to actually move the bytes. Data is moved one byte at a time.

In fig-FORTH:

```forth
addr1 addr2 n MOVE
```

Moves n sixteen-bit cells, beginning at addr1, to the memory locations beginning at addr2. The move proceeds from low to high memory. (From Forth Fundamentals Vol. 2, by Kevin McCabe.)

The time for test 1 was 7.2 seconds
The time for test 2 was 50.2 seconds.

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